Mathematical Tripos Part II: Michaelmas Term 2025

Numerical Analysis – Lecture 19

Algorithm 4.26 (The conjugate gradient method) Here it is.

- (A) For any initial vector $x^{(0)}$, set $d^{(0)} = r^{(0)} = b Ax^{(0)}$;
- (B) For $k \ge 0$, calculate $\boldsymbol{x}^{(k+1)} = \boldsymbol{x}^{(k)} + \alpha_k \boldsymbol{d}^{(k)}$ and the residual

$$\mathbf{r}^{(k+1)} = \mathbf{r}^{(k)} - \alpha_k A \mathbf{d}^{(k)}, \quad \text{with} \quad \alpha_k := \{ \mathbf{r}^{(k+1)} \perp \mathbf{d}^{(k)} \} = \frac{(\mathbf{r}^{(k)}, \mathbf{d}^{(k)})}{(A \mathbf{d}^{(k)}, \mathbf{d}^{(k)})}, \quad k \ge 0.$$
 (4.8)

(C) For the same k, the next conjugate direction is the vector

$$\mathbf{d}^{(k+1)} = \mathbf{r}^{(k+1)} + \beta_k \mathbf{d}^{(k)}, \quad \text{with} \quad \beta_k := \{ \mathbf{d}^{(k+1)} \perp A \mathbf{d}^{(k)} \} = -\frac{(\mathbf{r}^{(k+1)}, A \mathbf{d}^{(k)})}{(\mathbf{d}^{(k)}, A \mathbf{d}^{(k)})}, \quad k \ge 0. \quad (4.9)$$

Theorem 4.27 (Properties of CGM) For every $m \ge 0$, the conjugate gradient method has the following properties.

(1) The linear space spanned by the residuals $\{r^{(i)}\}$ is the same as the linear space spanned by the conjugate directions $\{d^{(i)}\}$ and it coincides with the space spanned by $\{A^ir^{(0)}\}$:

$$\mathrm{span}\{\boldsymbol{r}^{(i)}\}_{i=0}^{m} = \mathrm{span}\{\boldsymbol{d}^{(i)}\}_{i=0}^{m} = \mathrm{span}\{A^{i}\boldsymbol{r}^{(0)}\}_{i=0}^{m}.$$

- (2) The residuals satisfy the orthogonality conditions: $(\mathbf{r}^{(m)}, \mathbf{r}^{(i)}) = (\mathbf{r}^{(m)}, \mathbf{d}^{(i)}) = 0$ for i < m.
- (3) The directions are conjugate (A-orthogonal): $(\mathbf{d}^{(m)}, \mathbf{d}^{(i)})_A = (\mathbf{d}^{(m)}, A\mathbf{d}^{(i)}) = 0$ for i < m.

Proof. We use induction on $m \ge 0$, the assertions being trivial for m = 0, since $\mathbf{d}^{(0)} = \mathbf{r}^{(0)}$ and (2)-(3) are void. Therefore, assuming that the assertions are true for some m = k, we ask if they remain true when m = k + 1.

(1) Formula (4.9)

$$d^{(k+1)} = r^{(k+1)} + \beta_k d^{(k)}$$

readily implies that equivalence of the spaces spanned by $(\boldsymbol{r}^{(i)})_0^k$ and $(\boldsymbol{d}^{(i)})_0^k$, is preserved when k is increased to k+1. Similarly, from $\boldsymbol{r}^{(k+1)} = \boldsymbol{r}^{(k)} - \alpha_k \boldsymbol{A} \boldsymbol{d}^{(k)}$ in (4.8), and from the inductive assumption $\boldsymbol{r}^{(k)}, \boldsymbol{d}^{(k)} \in \operatorname{span}\{A^i \boldsymbol{r}^{(0)}\}_{i=0}^k$, it follows that $\boldsymbol{r}^{(k+1)} \in \operatorname{span}\{A^i \boldsymbol{r}^{(0)}\}_{i=0}^{k+1}$.

(2) Turning to assertion (2), we need $r^{(k+1)} \perp r^{(i)}$ for $i \leq k$, which by (1) is equivalent to

$$r^{(k+1)} \perp d^{(i)}$$
 for $i < k$.

We have $r^{(k+1)} \perp d^{(k)}$ by the definition of α_k in (4.8), so we need

$$m{r}^{(k+1)} \stackrel{(4.8)}{=} m{r}^{(k)} - lpha_k A m{d}^{(k)} \perp m{d}^{(i)} \quad ext{for} \quad i < k \,,$$

and this follow from the inductive assumptions $m{r}^{(k)} \perp m{d}^{(i)}$ and $Am{d}^{(k)} \perp m{d}^{(i)}$.

(3) It remains to justify (3), namely that $d^{(k+1)}$ defined in (4.9) satisfies

$$d^{(k+1)} \perp Ad^{(i)}$$
 for $i < k$.

The value of β_k in (4.9) is defined to give $d^{(k+1)} \perp Ad^{(k)}$, so we need

$$d^{(k+1)} \stackrel{(4.9)}{=} r^{(k+1)} + \beta_k d^{(k)} \perp A d^{(i)}$$
 for $i < k$.

By the inductive hypothesis $\boldsymbol{d}^{(k)} \perp A\boldsymbol{d}^{(i)}$, hence it remains to establish that $\boldsymbol{r}^{(k+1)} \perp A\boldsymbol{d}^{(i)}$ for i < k. Now, the formula (4.8) yields $A\boldsymbol{d}^{(i)} = (\boldsymbol{r}^{(i)} - \boldsymbol{r}^{(i+1)})/\alpha_i$, therefore we require the conditions $\boldsymbol{r}^{(k+1)} \perp (\boldsymbol{r}^{(i)} - \boldsymbol{r}^{(i+1)})$ for i < k, and they are a consequence of the assertion (2) for m = k+1 obtained previously.

Corollary 4.28 (A termination property) If the conjugate gradient method is applied in exact arithmetic, then, for any $\mathbf{x}^{(0)} \in \mathbb{R}^n$, termination occurs after at most n iterations. More precisely, termination occurs after at most s iterations, where $s = \dim \text{span}\{A^i \mathbf{r}_0\}_{i=0}^{n-1}$ (which can be smaller than n).

Proof. Assertion (2) of Theorem 4.27 states that residuals $(r^{(k)})_{k\geq 0}$ form a sequence of mutually orthogonal vectors in \mathbb{R}^n , therefore at most n of them can be nonzero. Since they also belong to the space $\operatorname{span}\{A^i r_0\}_{i=0}^{n-1}$, their number is bounded by the dimension of that space.

Definition 4.29 (The Krylov subspaces) Let A be an $n \times n$ matrix, $v \in \mathbb{R}^n$ nonzero, and $m \in \mathbb{N}$. The linear space $K_m(A, v) := \operatorname{span}\{A^i v\}_{i=0}^{m-1}$ is called the m-th Krylov subspace of \mathbb{R}^n .

Theorem 4.30 (Number of iterations in CGM) Let A > 0, and let s be the number of its distinct eigenvalues. Then, for any v,

$$\dim K_m(A, \mathbf{v}) \le s \quad \forall m. \tag{4.10}$$

Hence, for any A > 0, the number of iterations of the CGM for solving Ax = b is bounded by the number of distinct eigenvalues of A.

Proof. Inequality (4.10) is true not just for positive definite A>0, but for any A with n linearly independent eigenvectors (\boldsymbol{u}_i) . Indeed, in that case one can expand $\boldsymbol{v}=\sum_{i=1}^n a_i\boldsymbol{u}_i$, and then group together eigenvectors with the same eigenvalues: for each λ_{ν} we set $\boldsymbol{w}_{\nu}=\sum_{k=1}^{m_{\nu}}a_{i_k}\boldsymbol{u}_{i_k}$ if $A\boldsymbol{u}_{i_k}=\lambda_{\nu}\boldsymbol{u}_{i_k}$. Then

$$\boldsymbol{v} = \sum_{\nu=1}^{s} c_{\nu} \boldsymbol{w}_{\nu}, \qquad c_{\nu} \in \{0, 1\},$$

hence $A^i v = \sum_{\nu=1}^s c_\nu \lambda_\nu^i w_\nu$, thus for any m we get $K_m(A, v) \subseteq \operatorname{span}\{w_1, w_2, \dots, w_s\}$, and that proves (4.10). By Corollary 4.28, the number of iteration in CGM is bounded by $\dim K_m(A, r^{(0)})$, hence the final conclusion.

Remark 4.31 Theorem 4.30 shows that, unlike other iterative schemes, the conjugate gradient method is both iterative and direct: each iteration produces a reasonable approximation to the exact solution, and the exact solution itself will be recovered after n iterations at most.

We now simplify and reformulate Algorithm 4.26.

Firstly, we rewrite expressions for the parameters α_k and β_k in (4.8)-(4.9) as follows:

$$\alpha_{k} = \frac{(\boldsymbol{r}^{(k)}, \boldsymbol{d}^{(k)})}{(\boldsymbol{d}^{(k)}, A\boldsymbol{d}^{(k)})} \stackrel{(c)}{=} \frac{\|\boldsymbol{r}^{(k)}\|^{2}}{(\boldsymbol{d}^{(k)}, A\boldsymbol{d}^{(k)})} > 0,$$

$$\beta_{k} = -\frac{(\boldsymbol{r}^{(k+1)}, A\boldsymbol{d}^{(k)})}{(\boldsymbol{d}^{(k)}, A\boldsymbol{d}^{(k)})} \stackrel{(a)}{=} -\frac{(\boldsymbol{r}^{(k+1)}, \boldsymbol{r}^{(k+1)} - \boldsymbol{r}^{(k)})}{(\boldsymbol{d}^{(k)}, \boldsymbol{r}^{(k+1)} - \boldsymbol{r}^{(k)})} \stackrel{(b)}{=} \frac{\|\boldsymbol{r}^{(k+1)}\|^{2}}{(\boldsymbol{d}^{(k)}, \boldsymbol{r}^{(k)})} \stackrel{(c)}{=} \frac{\|\boldsymbol{r}^{(k+1)}\|^{2}}{\|\boldsymbol{r}^{(k)}\|^{2}} > 0.$$

Here, for β , we used in (a) the fact that $A\boldsymbol{d}^{(k)}$ is a multiple of $\boldsymbol{r}^{(k+1)}-\boldsymbol{r}^{(k)}$ by (4.8), and in (b) orthogonality of $\boldsymbol{r}^{(k+1)}$ to both $\boldsymbol{r}^{(k)},\boldsymbol{d}^{(k)}$ proved in Theorem 4.27(2). Then, for both β and α , we used in (c) the property $(\boldsymbol{d}^{(k)},\boldsymbol{r}^{(k)})=\|\boldsymbol{r}^{(k)}\|^2$ which follows from (4.9) with index k+1, taking in account orthogonality $\boldsymbol{r}^{(k+1)}\perp\boldsymbol{d}^{(k)}$.

Secondly, we let $\hat{x}^{(0)}$ be the zero vector.

Algorithm 4.32 (Standard form of the conjugate gradient method) Here it is.

- (1) Set k = 0, $\mathbf{x}^{(0)} = 0$, $\mathbf{r}^{(0)} = \mathbf{b}$, and $\mathbf{d}^{(0)} = \mathbf{r}^{(0)}$;
- (2) Calculate the matrix-vector product $v^{(k)} = Ad^{(k)}$ and $\alpha_k = ||r^{(k)}||^2/(d^{(k)}, v^{(k)}) > 0$;
- (3) Apply the formulae $x^{(k+1)} = x^{(k)} + \alpha_k d^{(k)}$ and $r^{(k+1)} = r^{(k)} \alpha_k v^{(k)}$;
- (4) Stop if $||r^{(k+1)}||$ is acceptably small;
- (5) Set $d^{(k+1)} = r^{(k+1)} + \beta_k d^{(k)}$, where $\beta_k = ||r^{(k+1)}||^2 / ||r^{(k)}||^2 > 0$;
- (6) Increase $k \rightarrow k + 1$ and go back to (2).

The total work is dominated by the number of iterations, multiplied by the time it takes to compute $v^{(k)} = Ad^{(k)}$. Thus the conjugate gradient algorithm is highly suitable when most of the elements of A are zero, i.e. when A is *sparse*.