

Mathematical Tripod Part II: Michaelmas Term 2021

Numerical Analysis – Lecture 19

Conjugate gradient method The conjugate gradient method is the method of conjugate directions (Theorem 4.23 from previous lecture) where the directions $\mathbf{d}^{(i)}$ are chosen so that they A -orthogonalize the residuals, i.e., the $\mathbf{d}^{(i)}$ satisfy

$$\text{span}(\mathbf{d}^{(0)}, \dots, \mathbf{d}^{(k-1)}) = \text{span}(\mathbf{r}^{(0)}, \dots, \mathbf{r}^{(k-1)}) \quad (4.8)$$

for every iteration k , in addition to being pairwise A -orthogonal. This can be achieved by setting $\mathbf{d}^{(0)} = \mathbf{r}^{(0)}$ and applying the Gram-Schmidt step at each iteration

$$\mathbf{d}^{(k+1)} = \mathbf{r}^{(k+1)} - \sum_{i \leq k} \frac{\langle \mathbf{r}^{(k+1)}, \mathbf{d}^{(i)} \rangle_A}{\langle \mathbf{d}^{(i)}, \mathbf{d}^{(i)} \rangle_A} \mathbf{d}^{(i)}. \quad (4.9)$$

Because of our particular choice of $\mathbf{d}^{(k)}$, the equation above simplifies dramatically and the terms $i \leq k-1$ in the summation above happen to be zero! This is the *key point* of the CG method. Let's prove this. Recall that the iterates are defined by $\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \alpha_k \mathbf{d}^{(k)}$ so that the residuals satisfy $\mathbf{r}^{(k+1)} = \mathbf{r}^{(k)} - \alpha_k A \mathbf{d}^{(k)}$. It is easy then to see by induction, using the property (4.8) that

$$\text{span}\{\mathbf{d}^{(i)}\}_{i=0}^{k-1} = \text{span}\{\mathbf{r}^{(i)}\}_{i=0}^{k-1} = \text{span}\{\mathbf{r}^{(0)}, A\mathbf{r}^{(0)}, \dots, A^{k-1}\mathbf{r}^{(0)}\} =: K_k(A, \mathbf{r}^{(0)}),$$

where $K_m(A, \mathbf{v}) = \text{span}\{A^i \mathbf{v}\}_{i=0}^{m-1}$ is the m 'th Krylov subspace of A wrt \mathbf{v} . The result of Theorem 4.23 tells us that $\mathbf{r}^{(k+1)}$ is orthogonal to $K_{k+1}(A, \mathbf{r}^{(0)})$. Now for $i < k$, we have $\mathbf{d}^{(i)} \in K_k(A, \mathbf{r}^{(0)})$ and so $A\mathbf{d}^{(i)} \in K_{k+1}(A, \mathbf{r}^{(0)})$. This implies that $\langle \mathbf{r}^{(k+1)}, A\mathbf{d}^{(i)} \rangle = 0$ for $i < k$, and shows that the terms $i < k$ in Equation (4.9) are equal to zero.

The conjugate gradient algorithm can thus be summarized in the following: Set $\mathbf{d}^{(0)} = \mathbf{r}^{(0)} = \mathbf{b} - A\mathbf{x}^{(0)}$ and iterate, for $k \geq 0$:

$$\begin{cases} \mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \alpha_k \mathbf{d}^{(k)} & \alpha_k = \frac{\langle \mathbf{r}^{(k)}, \mathbf{d}^{(k)} \rangle}{\langle \mathbf{d}^{(k)}, A\mathbf{d}^{(k)} \rangle} \\ \mathbf{d}^{(k+1)} = \mathbf{r}^{(k+1)} + \beta_k \mathbf{d}^{(k)} & \beta_k = -\frac{\langle \mathbf{r}^{(k+1)}, A\mathbf{d}^{(k)} \rangle}{\langle \mathbf{d}^{(k)}, A\mathbf{d}^{(k)} \rangle} \end{cases} \quad (4.10)$$

where $\mathbf{r}^{(k)}$ stands for $\mathbf{b} - A\mathbf{x}^{(k)}$. We can summarize the properties of the Conjugate Gradient Method in the following theorem.

Theorem 4.26 (Properties of CGM) For every $m \geq 0$, the conjugate gradient method has the following properties.

- (1) The linear space spanned by the residuals $\{\mathbf{r}^{(i)}\}$ is the same as the linear space spanned by the conjugate directions $\{\mathbf{d}^{(i)}\}$ and it coincides with the space spanned by $\{A^i \mathbf{r}^{(0)}\}$:

$$\text{span}\{\mathbf{r}^{(i)}\}_{i=0}^m = \text{span}\{\mathbf{d}^{(i)}\}_{i=0}^m = \text{span}\{A^i \mathbf{r}^{(0)}\}_{i=0}^m.$$

- (2) The residuals satisfy the orthogonality conditions: $\langle \mathbf{r}^{(m)}, \mathbf{r}^{(i)} \rangle = \langle \mathbf{r}^{(m)}, \mathbf{d}^{(i)} \rangle = 0$ for $i < m$.
- (3) The directions are conjugate (A -orthogonal): $\langle \mathbf{d}^{(m)}, \mathbf{d}^{(i)} \rangle_A = \langle \mathbf{d}^{(m)}, A\mathbf{d}^{(i)} \rangle = 0$ for $i < m$.

Using these properties we can simplify the expressions for α_k and β_k . Indeed, using the second equation in (4.10), and the fact that $\mathbf{r}^{(k)} \perp \mathbf{d}^{(k-1)}$, we have

$$\langle \mathbf{r}^{(k)}, \mathbf{d}^{(k)} \rangle = \langle \mathbf{r}^{(k)}, \mathbf{r}^{(k)} \rangle = \|\mathbf{r}^{(k)}\|_2^2 \quad (4.11)$$

which shows that

$$\alpha_k = \frac{\|\mathbf{r}^{(k)}\|_2^2}{\langle \mathbf{d}^{(k)}, A\mathbf{d}^{(k)} \rangle} > 0.$$

Also, we can write:

$$\beta_k = -\frac{\langle \mathbf{r}^{(k+1)}, A\mathbf{d}^{(k)} \rangle}{\langle \mathbf{d}^{(k)}, A\mathbf{d}^{(k)} \rangle} \stackrel{(a)}{=} -\frac{\langle \mathbf{r}^{(k+1)}, \mathbf{r}^{(k+1)} - \mathbf{r}^{(k)} \rangle}{\langle \mathbf{d}^{(k)}, \mathbf{r}^{(k+1)} - \mathbf{r}^{(k)} \rangle} \stackrel{(b)}{=} \frac{\|\mathbf{r}^{(k+1)}\|^2}{\langle \mathbf{d}^{(k)}, \mathbf{r}^{(k)} \rangle} \stackrel{(c)}{=} \frac{\|\mathbf{r}^{(k+1)}\|^2}{\|\mathbf{r}^{(k)}\|^2} > 0.$$

where we used in (a) the fact that $A\mathbf{d}^{(k)}$ is a multiple of $\mathbf{r}^{(k+1)} - \mathbf{r}^{(k)}$, and in (b) orthogonality of $\mathbf{r}^{(k+1)}$ to both $\mathbf{r}^{(k)}$, $\mathbf{d}^{(k)}$ (Theorem 4.26(2)), and in (c) we used (4.11).

The complete conjugate gradient method can thus be written as follows:

Algorithm 4.27 (Standard form of the conjugate gradient method) –

- (1) Set $k = 0$, $\mathbf{x}^{(0)} = 0$, $\mathbf{r}^{(0)} = \mathbf{b}$, and $\mathbf{d}^{(0)} = \mathbf{r}^{(0)}$;
- (2) Calculate the matrix-vector product $\mathbf{v}^{(k)} = A\mathbf{d}^{(k)}$ and $\alpha_k = \|\mathbf{r}^{(k)}\|^2 / \langle \mathbf{d}^{(k)}, \mathbf{v}^{(k)} \rangle > 0$;
- (3) Apply the formulae $\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \alpha_k \mathbf{d}^{(k)}$ and $\mathbf{r}^{(k+1)} = \mathbf{r}^{(k)} - \alpha_k \mathbf{v}^{(k)}$;
- (4) Stop if $\|\mathbf{r}^{(k+1)}\|$ is acceptably small;
- (5) Set $\mathbf{d}^{(k+1)} = \mathbf{r}^{(k+1)} + \beta_k \mathbf{d}^{(k)}$, where $\beta_k = \|\mathbf{r}^{(k+1)}\|^2 / \|\mathbf{r}^{(k)}\|^2 > 0$;
- (6) Increase $k \rightarrow k + 1$ and go back to (2).

The total work is dominated by the number of iterations, multiplied by the time it takes to compute $\mathbf{v}^{(k)} = A\mathbf{d}^{(k)}$. Thus the conjugate gradient algorithm is highly suitable when most of the elements of A are zero, i.e. when A is *sparse*.

Finite termination We have already seen that the method of conjugate directions (Theorem 4.23 in previous lecture) terminates after at most n steps. We restate this result in the special case of the conjugate gradient method.

Corollary 4.28 (A termination property) *If the conjugate gradient method is applied in exact arithmetic, then, for any $\mathbf{x}^{(0)} \in \mathbb{R}^n$, termination occurs after at most n iterations. More precisely, termination occurs after at most s iterations, where $s = \dim \text{span}\{A^i \mathbf{r}_0\}_{i=0}^{n-1}$ (which can be smaller than n).*

Proof. Assertion (2) of Theorem 4.26 states that residuals $(\mathbf{r}^{(k)})_{k \geq 0}$ form a sequence of mutually orthogonal vectors in \mathbb{R}^n , therefore at most n of them can be nonzero. Since they also belong to the space $\text{span}\{A^i \mathbf{r}_0\}_{i=0}^{n-1}$, their number is bounded by the dimension of that space. \square

We can bound the dimension of the Krylov subspace $\text{span}\{A^i \mathbf{r}_0\}_{i=0}^{n-1}$ using the number of distinct eigenvalues of A .

Theorem 4.29 (Number of iterations in CGM) *Let $A > 0$, and let s be the number of its distinct eigenvalues. Then, for any \mathbf{v} ,*

$$\dim K_m(A, \mathbf{v}) \leq s \quad \forall m. \quad (4.12)$$

Hence, for any $A > 0$, the number of iterations of the CGM for solving $A\mathbf{x} = \mathbf{b}$ is bounded by the number of distinct eigenvalues of A .

Proof. Inequality (4.12) is true not just for positive definite $A > 0$, but for any A with n linearly independent eigenvectors (\mathbf{u}_i) . Indeed, in that case one can expand $\mathbf{v} = \sum_{i=1}^n a_i \mathbf{u}_i$, and then group together eigenvectors with the same eigenvalues: for each λ_ν we set $\mathbf{w}_\nu = \sum_{k=1}^{m_\nu} a_{i_k} \mathbf{u}_{i_k}$ if $A\mathbf{u}_{i_k} = \lambda_\nu \mathbf{u}_{i_k}$. Then

$$\mathbf{v} = \sum_{\nu=1}^s c_\nu \mathbf{w}_\nu, \quad c_\nu \in \{0, 1\},$$

hence $A^i \mathbf{v} = \sum_{\nu=1}^s c_\nu \lambda_\nu^i \mathbf{w}_\nu$, thus for any m we get $K_m(A, \mathbf{v}) \subseteq \text{span}\{\mathbf{w}_1, \mathbf{w}_2, \dots, \mathbf{w}_s\}$, and that proves (4.12). By Corollary 4.28, the number of iteration in CGM is bounded by $\dim K_m(A, \mathbf{r}^{(0)})$, hence the final conclusion. \square

Remark 4.30 Theorem 4.29 shows that, unlike other iterative schemes, the conjugate gradient method is both iterative and direct: each iteration produces a reasonable approximation to the exact solution, and the exact solution itself will be recovered after n iterations at most.

Convergence One can prove a more quantitative version of Theorem (4.29).

Theorem 4.31 *Let A be symmetric positive definite. After k iterations of the conjugate gradient method, the error $e^{(k)} = \mathbf{x}^* - \mathbf{x}^{(k)}$ satisfies*

$$\|e^{(k)}\|_A = \min_{P_k} \|P_k(A)e^{(0)}\|_A$$

where the minimization is over all polynomials P_k of degree $\leq k$ that satisfy $P_k(0) = 1$.

Proof. We know from Lecture 18 (Equation (4.7) applied recursively) that $e^{(k)}$ is obtained from $e^{(0)}$ by projecting out (in the inner product $\langle \cdot, \cdot \rangle_A$) the components $\mathbf{d}^{(0)}, \dots, \mathbf{d}^{(k-1)}$. This means that

$$\|e^{(k)}\|_A = \min_{\mathbf{v}} \|e^{(0)} - \mathbf{v}\|_A$$

where the minimization is over all $\mathbf{v} \in \text{span}(\mathbf{d}^{(0)}, \dots, \mathbf{d}^{(k-1)})$. For the conjugate gradient method, this subspace is the same as $\text{span}(\mathbf{r}^{(0)}, \dots, A^{k-1}\mathbf{r}^{(0)})$, and since $\mathbf{r}^{(0)} = Ae^{(0)}$, this means that any such \mathbf{v} can be written as $\mathbf{v} = \sum_{i=1}^k c_i A^i e^{(0)}$. Let $P_k(t) = 1 - \sum_{i=1}^k c_i t^i$ we get the desired equality.

Remark 4.32 *If A has s distinct eigenvalues $\lambda_1, \dots, \lambda_s > 0$, then with $P_s(t) = \prod_{i=1}^s (1 - t/\lambda_i)$ we have $\deg P_s = s$, $P_s(0) = 1$, and $P_s(A) = 0$. Thus this shows that the CG method terminates after s iterations.*