

Mathematical Tripos Part II: Michaelmas Term 2022

Numerical Analysis – Lecture 4

Algorithm 1.15 (The fast Fourier transform (FFT)) We assume that n is a power of 2, i.e. $n = 2m = 2^p$, and for $\mathbf{y} \in \mathbb{C}^{2m}$, denote by

$$\mathbf{y}^{(E)} = \{y_{2j}\}_{j=0,\dots,m} \in \mathbb{C}^m \quad \text{and} \quad \mathbf{y}^{(O)} = \{y_{2j+1}\}_{j=0,\dots,m} \in \mathbb{C}^m$$

the even and odd portions of \mathbf{y} , respectively.

Suppose that we already know the inverse DFT of both ‘short’ sequences,

$$\mathbf{x}^{(E)} = \mathcal{F}_m \mathbf{y}^{(E)}, \quad \mathbf{x}^{(O)} = \mathcal{F}_m \mathbf{y}^{(O)}.$$

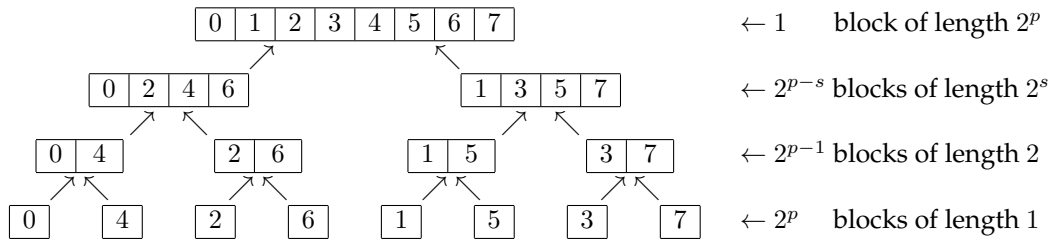
It is then possible to assemble $\mathbf{x} = \mathcal{F}_{2m} \mathbf{y}$ in a small number of operations. Indeed, for $\ell \in \{0, \dots, m-1\}$, we have

$$\begin{aligned} x_\ell &= \sum_{j=0}^{2m-1} \omega_{2m}^{j\ell} y_j = \sum_{j=0}^{m-1} \omega_{2m}^{2j\ell} y_{2j} + \sum_{j=0}^{m-1} \omega_{2m}^{(2j+1)\ell} y_{2j+1} \\ &= \sum_{j=0}^{m-1} \omega_m^{j\ell} y_j^{(E)} + \omega_{2m}^\ell \sum_{j=0}^{m-1} \omega_m^{j\ell} y_j^{(O)} = x_\ell^{(E)} + \omega_{2m}^\ell x_\ell^{(O)}. \end{aligned}$$

Therefore, it costs just m products to evaluate the first half of \mathbf{x} , provided that $\mathbf{x}^{(E)}$ and $\mathbf{x}^{(O)}$ are known. It actually costs nothing to evaluate the second half, since

$$\omega_m^{j(m+\ell)} = \omega_m^{j\ell}, \quad \omega_{2m}^{m+\ell} = -\omega_{2m}^\ell \Rightarrow x_{m+\ell} = x_\ell^{(E)} - \omega_{2m}^\ell x_\ell^{(O)}, \quad \ell = 0, \dots, m-1.$$

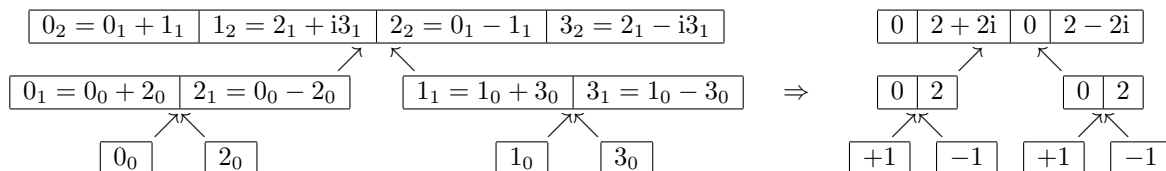
To execute FFT, we start from vectors of unit length and in each s -th stage, $s = 1 \dots p$, assemble 2^{p-s} vectors of length 2^s from vectors of length 2^{s-1} : this costs $2^{p-s} 2^{s-1} = 2^{p-1}$ products. Altogether, the cost of FFT is $p 2^{p-1} = \frac{1}{2} n \log_2 n$ products.



For $n = 1024 = 2^{10}$, say, the cost is $\approx 5 \times 10^3$ products, compared to $\approx 10^6$ for naive matrix multiplication! For $n = 2^{20}$ the respective numbers are $\approx 1.05 \times 10^7$ and $\approx 1.1 \times 10^{12}$, which represents a saving by a factor of more than 10^5 .

Matlab demo: Check out the online animation for computing the FFT at http://www.damtp.cam.ac.uk/user/hf323/M21-II-NA/demos/fft_gui/fft_gui.html and download the Matlab GUI from there to follow the computation of each single FFT term.

Example 1.16 Computation of FFT for $n = 4$ in general, and for the vector $\mathbf{y} = (1, 1, -1, -1)$ in particular.



2 Partial differential equations of evolution

We consider the *diffusion equation*

$$\frac{\partial u}{\partial t} = \frac{\partial^2 u}{\partial x^2}, \quad 0 \leq x \leq 1, \quad t \geq 0,$$

with *initial conditions* $u(x, 0) = u_0(x)$ for $t = 0$ and *Dirichlet boundary conditions* $u(0, t) = \phi_0(t)$ at $x = 0$ and $u(1, t) = \phi_1(t)$ at $x = 1$. By Taylor's expansion

$$\begin{aligned} \frac{\partial u(x, t)}{\partial t} &= \frac{1}{k} [u(x, t+k) - u(x, t)] + \mathcal{O}(k), & k = \Delta t, \\ \frac{\partial^2 u(x, t)}{\partial x^2} &= \frac{1}{h^2} [u(x-h, t) - 2u(x, t) + u(x+h, t)] + \mathcal{O}(h^2), & h = \Delta x, \end{aligned}$$

so that, for the exact solution $u = \hat{u}$ of the diffusion equation, we obtain

$$\hat{u}(x, t+k) = \hat{u}(x, t) + \frac{k}{h^2} [\hat{u}(x-h, t) - 2\hat{u}(x, t) + \hat{u}(x+h, t)] + \eta(x, t) \quad (2.1)$$

where $\eta(x, t) = \mathcal{O}(k^2 + kh^2)$. (More precisely, one proves using Taylor's theorem that $|\eta(x, t)| \leq c_1 k^2 + c_2 kh^2$ where $c_1 = \frac{1}{2} \max_{\xi, \tau} |\frac{\partial^2 \hat{u}}{\partial t^2}(\xi, \tau)|$ and $c_2 = \frac{1}{12} \max_{\xi, \tau} |\frac{\partial^4 \hat{u}}{\partial x^4}(\xi, \tau)|$.) That motivates the numerical scheme for approximation $u_m^n \approx \hat{u}(x_m, t_n)$ on the rectangular mesh $(x_m, t_n) = (mh, nk)$:

$$u_m^{n+1} = u_m^n + \mu (u_{m-1}^n - 2u_m^n + u_{m+1}^n), \quad m = 1 \dots M. \quad (2.2)$$

Here $h = \frac{1}{M+1}$ and $\mu = \frac{k}{h^2} = \frac{\Delta t}{(\Delta x)^2}$ is the so-called *Courant number*. With μ being fixed, we have $k = \mu h^2$, so that the local truncation error of the scheme is $\mathcal{O}(h^4)$. Substituting whenever necessary initial conditions u_m^0 and boundary conditions u_0^n and u_{M+1}^n , we possess enough information to advance in (2.2) from $\mathbf{u}^n := [u_1^n, \dots, u_M^n]$ to $\mathbf{u}^{n+1} := [u_1^{n+1}, \dots, u_M^{n+1}]$.

Similarly to ODEs or Poisson equation, we say that the method is *convergent* if, for a fixed μ , and for every $T > 0$, we have

$$\lim_{\substack{h \rightarrow 0, k \rightarrow 0 \\ k/h^2 = \mu}} \max_{\substack{1 \leq m \leq M \\ 1 \leq n \leq T/k}} |u_m^n - \hat{u}(mh, nk)| = 0.$$

Theorem 2.1 *If $\mu \leq \frac{1}{2}$, then method (2.2) converges.*

Proof. Let $e_m^n := \hat{u}(mh, nk) - u_m^n$ be the error of approximation, and let $\mathbf{e}^n = [e_1^n, \dots, e_M^n]$ with $\|\mathbf{e}^n\|_\infty := \max_m |e_m^n|$. Convergence is equivalent to

$$\lim_{h \rightarrow 0} \max_{1 \leq n \leq T/k} \|\mathbf{e}^n\|_\infty = 0$$

for every constant $T > 0$. Subtracting (2.1) from (2.2), we obtain

$$\begin{aligned} e_m^{n+1} &= e_m^n + \mu (e_{m-1}^n - 2e_m^n + e_{m+1}^n) + \eta_m^n \\ &= \mu e_{m-1}^n + (1 - 2\mu) e_m^n + \mu e_{m+1}^n + \eta_m^n \end{aligned}$$

where $|\eta_m^n| \leq ch^4$ for some constant $c > 0$ (namely $c = c_1 \mu^2 + c_2 \mu$, where $c_1, c_2 > 0$ are defined after equation (2.1)). Then

$$\|\mathbf{e}^{n+1}\|_\infty = \max_m |e_m^{n+1}| \leq (2\mu + |1 - 2\mu|) \|\mathbf{e}^n\|_\infty + ch^4 = \|\mathbf{e}^n\|_\infty + ch^4,$$

by virtue of $\mu \leq \frac{1}{2}$. Since $\|\mathbf{e}^0\|_\infty = 0$, induction yields

$$\|\mathbf{e}^n\|_\infty \leq cnh^4 \leq \frac{cT}{k} h^4 = \frac{cT}{\mu} h^2 \rightarrow 0 \quad (h \rightarrow 0) \quad \square$$

Matlab demo: Download the Matlab GUI for *Stability of 1D PDEs* from http://www.damtp.cam.ac.uk/user/hf323/M21-II-NA/demos/pde_stability/pde_stability.html and solve the diffusion equation in the interval $[0, 1]$ with method (2.2) and $\mu = 0.51 > \frac{1}{2}$. Using (as preset) 100 grid points to discretise $[0, 1]$ will then require the time steps to be $5.1 \cdot 10^{-5}$. The solution will evolve very slowly, but wait long enough to see what happens!