Mathematical Tripos Part II: Michaelmas Term 2023 Numerical Analysis – Lecture 4

Fast Fourier Transform (FFT) We assume that *n* is a power of 2, i.e. $n = 2m = 2^p$, and for $y \in \mathbb{C}^{2m}$, denote by

$$\boldsymbol{y}^{(\mathrm{E})} = \{y_{2j}\}_{j=0,...,m} \in \mathbb{C}^m \quad \text{and} \quad \boldsymbol{y}^{(\mathrm{O})} = \{y_{2j+1}\}_{j=0,...,m} \in \mathbb{C}^m$$

the even and odd portions of y, respectively.

Suppose that we already know the DFT of both 'short' sequences,

$$\boldsymbol{x}^{(\mathrm{E})} = \mathcal{F}_m \boldsymbol{y}^{(\mathrm{E})}, \qquad \boldsymbol{x}^{(\mathrm{O})} = \mathcal{F}_m \boldsymbol{y}^{(\mathrm{O})}.$$

It is then possible to assemble $x = \mathcal{F}_{2m}y$ in a small number of operations. Indeed, for $\ell \in \{0, ..., m-1\}$, we have

$$\begin{aligned} x_{\ell} &= \sum_{j=0}^{2m-1} \omega_{2m}^{j\ell} y_{j} &= \sum_{j=0}^{m-1} \omega_{2m}^{2j\ell} y_{2j} + \sum_{j=0}^{m-1} \omega_{2m}^{(2j+1)\ell} y_{2j+1} \\ &= \sum_{j=0}^{m-1} \omega_{m}^{j\ell} y_{j}^{(\mathrm{E})} + \omega_{2m}^{\ell} \sum_{j=0}^{m-1} \omega_{m}^{j\ell} y_{j}^{(\mathrm{O})} = x_{\ell}^{(\mathrm{E})} + \omega_{2m}^{\ell} x_{\ell}^{(\mathrm{O})}. \end{aligned}$$

Therefore, it costs just *m* products to evaluate the first half of *x*, provided that $x^{(E)}$ and $x^{(O)}$ are known. It actually costs nothing to evaluate the second half, since

$$\omega_m^{j(m+\ell)} = \omega_m^{j\ell}, \qquad \omega_{2m}^{m+\ell} = -\omega_{2m}^{\ell} \qquad \Rightarrow \qquad x_{m+\ell} = x_\ell^{(\mathrm{E})} - \omega_{2m}^{\ell} x_\ell^{(\mathrm{O})}, \qquad \ell = 0, \dots, m-1.$$

To execute FFT, we start from vectors of unit length and in each *s*-th stage, s = 1...p, assemble 2^{p-s} vectors of length 2^s from vectors of length 2^{s-1} : this costs $2^{p-s}2^{s-1} = 2^{p-1}$ products. Altogether, the cost of FFT is $p2^{p-1} = \frac{1}{2}n \log_2 n$ products.



For $n = 1024 = 2^{10}$, say, the cost is $\approx 5 \times 10^3$ products, compared to $\approx 10^6$ for naive matrix multiplication! For $n = 2^{20}$ the respective numbers are $\approx 1.05 \times 10^7$ and $\approx 1.1 \times 10^{12}$, which represents a saving by a factor of more than 10^5 .

Matlab demo: Check out the online animation for computing the FFT at http://www.damtp.cam.ac. uk/user/hf323/M21-II-NA/demos/fft_gui/fft_gui.html and download the Matlab GUI from there to follow the computation of each single FFT term.

Example 1.15 Computation of FFT for n = 4 in general, and for the vector $\boldsymbol{y} = (1, 1, -1, -1)$ in particular.



2 Partial differential equations of evolution

We consider the diffusion equation

$$\frac{\partial u}{\partial t} = \frac{\partial^2 u}{\partial x^2}, \qquad 0 \le x \le 1, \quad t \ge 0,$$

with *initial conditions* $u(x, 0) = u_0(x)$ for t = 0 and zero Dirichlet boundary conditions u(0, t) = u(1, t) = 0. By Taylor's expansion

$$\begin{array}{rcl} \frac{\partial u(x,t)}{\partial t} &=& \frac{1}{k} \big[u(x,t+k) - u(x,t) \big] + \mathcal{O}(k), & k = \Delta t \,, \\ \frac{\partial^2 u(x,t)}{\partial x^2} &=& \frac{1}{h^2} \big[u(x-h,t) - 2u(x,t) + u(x+h,t) \big] + \mathcal{O}(h^2), & h = \Delta x \,, \end{array}$$

so that, for the exact solution $u = \hat{u}$ of the diffusion equation, we obtain

$$\widehat{u}(x,t+k) = \widehat{u}(x,t) + \frac{k}{h^2} \left[\widehat{u}(x-h,t) - 2\widehat{u}(x,t) + \widehat{u}(x+h,t) \right] + \eta(x,t)$$
(2.1)

where $\eta(x,t) = \mathcal{O}(k^2 + kh^2)$. (More precisely, one proves using Taylor's theorem that $|\eta(x,t)| \leq c_1 k^2 + c_2 kh^2$ where $c_1 = \frac{1}{2} \max_{\xi,\tau} |\frac{\partial^2 \hat{u}}{\partial t^2}(\xi,\tau)|$ and $c_2 = \frac{1}{12} \max_{\xi,\tau} |\frac{\partial^4 \hat{u}}{\partial x^4}(\xi,\tau)|$.) That motivates the numerical scheme for approximation $u_m^n \approx \hat{u}(x_m, t_n)$ on the rectangular mesh $(x_m, t_n) = (mh, nk)$:

$$u_m^{n+1} = u_m^n + \mu \left(u_{m-1}^n - 2u_m^n + u_{m+1}^n \right), \qquad m = 1...M.$$
(2.2)

Here $h = \frac{1}{M+1}$ and $\mu = \frac{k}{h^2} = \frac{\Delta t}{(\Delta x)^2}$ is the so-called *Courant number*. With μ being fixed, we have $k = \mu h^2$, so that the local truncation error of the scheme is $\mathcal{O}(k^2)$. Substituting whenever necessary initial conditions u_m^0 and boundary conditions u_0^n and u_{M+1}^n , we possess enough information to advance in (2.2) from $u^n := [u_1^n, \ldots, u_M^n]$ to $u^{n+1} := [u_1^{n+1}, \ldots, u_M^{n+1}]$.

Similarly to ODEs or Poisson equation, we say that the method is *convergent* if, for a fixed μ , and for every T > 0, we have

$$\lim_{\substack{h \to 0, k \to 0 \\ k/h^2 = \mu}} \max_{\substack{1 \le m \le M \\ 1 \le n \le T/k}} |u_m^n - \widehat{u}(mh, nk)| = 0.$$

Theorem 2.1 If $\mu \leq \frac{1}{2}$, then method (2.2) converges.

Proof. Let $e_m^n := \hat{u}(mh, nk) - u_m^n$ be the error of approximation, and let $e^n = [e_1^n, \dots, e_M^n]$ with $||e^n||_{\infty} := \max_m |e_m^n|$. Convergence is equivalent to

$$\lim_{h \to 0} \max_{1 \le n \le T/k} \|\boldsymbol{e}^n\|_{\infty} = 0$$

for every constant T > 0. Subtracting (2.1) from (2.2), we obtain

$$e_m^{n+1} = e_m^n + \mu(e_{m-1}^n - 2e_m^n + e_{m+1}^n) + \eta_m^n$$

= $\mu e_{m-1}^n + (1 - 2\mu)e_m^n + \mu e_{m+1}^n + \eta_m^n$

where $|\eta_m^n| \le ck^2$ for some constant c > 0 (namely $c = c_1 + c_2/\mu$, where $c_1, c_2 > 0$ are defined after equation (2.1)). Then

$$\|e^{n+1}\|_{\infty} = \max_{m} |e_{m}^{n+1}| \le (2\mu + |1 - 2\mu|) \|e^{n}\|_{\infty} + ck^{2} = \|e^{n}\|_{\infty} + ck^{2},$$

by virtue of $\mu \leq \frac{1}{2}$. Since $\|e^0\|_{\infty} = 0$, induction yields

$$\|\boldsymbol{e}^n\|_{\infty} \le cnk^2 \le \frac{cT}{k}k^2 = cTk \to 0 \qquad (k \to 0).$$

Matlab demo: Download the Matlab GUI for *Stability of 1D PDEs* from http://www.damtp.cam.ac. uk/user/hf323/M21-II-NA/demos/pde_stability/pde_stability.html and solve the diffusion equation in the interval [0,1] with method (2.2) and $\mu = 0.51 > \frac{1}{2}$. Using (as preset) 100 grid points to discretise [0,1] will then require the time steps to be $5.1 \cdot 10^{-5}$. The solution will evolve very slowly, but wait long enough to see what happens!