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**Exercises for revision class****Contents**

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## 1 Diagonally dominant matrices

A matrix  $A \in \mathbf{S}^n$  is called *diagonally dominant* if  $A_{ii} \geq \sum_{j \neq i} |A_{ij}|$  for all  $i = 1, \dots, n$ . Let  $\mathcal{D}_n \subset \mathbf{S}^n$  be the set of diagonally dominant matrices.

- (a) Show that if  $A$  is diagonally dominant then it is positive semidefinite.
- (b) Recall the definition of *proper cone*. Show that the set  $\mathcal{D}_n$  is a proper cone in  $\mathbf{S}^n$ .
- (c) Show that the extreme rays of  $\mathcal{D}_n$  are spanned by the matrices

$$e_i e_i^T \quad (i = 1, \dots, n) \quad \text{and} \quad (e_i \pm e_j)(e_i \pm e_j)^T \quad (1 \leq i < j \leq n).$$

where  $e_i \in \mathbb{R}^n$  is the vector with 1 in the  $i$ 'th component and 0 elsewhere.

## 2 Euclidean distance matrices

Let  $(d_{ij})_{1 \leq i < j \leq n}$  be  $\binom{n}{2}$  positive numbers. Show that the following two assertions are equivalent:

- (i) There exist points  $x_1, \dots, x_n \in \mathbb{R}^k$  (for some  $k$ ) such that  $d_{ij} = \|x_i - x_j\|_2$  for all  $i < j$ .
- (ii) The  $n \times n$  symmetric matrix  $D = [d_{ij}^2]_{1 \leq i, j \leq n}$  (where  $d_{ii} = 0$ ) is negative semidefinite on the subspace orthogonal to  $e = (1, \dots, 1) \in \mathbb{R}^n$ . [We say that a matrix  $A \in \mathbf{S}^n$  is *negative semidefinite on a subspace  $L$*  if  $x^T A x \leq 0$  for all  $x \in L$ ]

## 3 Binary quadratic optimisation: Nesterov's $2/\pi$ result

Let  $A$  be a real symmetric matrix of size  $n \times n$ , and consider the following binary quadratic optimisation problem:

$$\text{maximise } x^T A x \quad : \quad x \in \{-1, 1\}^n. \tag{1}$$

Let  $v^*$  be the optimal value of (1).

- (a) Consider the semidefinite program:

$$\text{maximise } \text{Tr}(AX) \quad : \quad X \succeq 0 \text{ and } X_{ii} = 1, \forall i = 1, \dots, n. \tag{2}$$

Let  $p_{SDP}^*$  be the optimal value of (2). Show that  $v^* \leq p_{SDP}^*$ .

**From now on we are going to assume that  $A$  is positive semidefinite.** The purpose of the rest of this problem is to show that  $\frac{2}{\pi} p_{SDP}^* \leq v^*$ . To prove this inequality, we will use a “randomised rounding” scheme similar to the one we saw in lecture for the maximum cut problem.

- (b) Let  $X$  be the optimal solution (2) and let  $v_1, \dots, v_n \in \mathbb{R}^r$  with  $r = \text{rank}(X)$  such that  $X_{ij} = \langle v_i, v_j \rangle$  for all  $i, j = 1, \dots, n$ . Define the random variable  $y \in \{-1, 1\}^n$  as follows:

$$y_i = \text{sign}(\langle v_i, Z \rangle)$$

where  $Z$  is a standard Gaussian variable on  $\mathbb{R}^r$ . We saw in lecture that

$$\mathbb{E}[y_i y_j] = \frac{2}{\pi} \arcsin(X_{ij}) \quad \forall 1 \leq i, j \leq n,$$

which you can use without proof. Show that:

$$v^* \geq \mathbb{E}[y^T A y] = \frac{2}{\pi} \text{Tr}(A \arcsin[X]).$$

where  $\arcsin[X]$  is the matrix obtained by applying the arcsin function to each entry of  $X$ , i.e.,  $\arcsin[X]_{ij} = \arcsin(X_{ij})$ .

(c) Recall the Schur product theorem:

*Schur product theorem:* If  $P \succeq 0$  and  $Q \succeq 0$  then  $P \odot Q \succeq 0$  where  $P \odot Q$  is the entrywise product of  $P$  and  $Q$ .

Use the Schur product theorem (without proof) to show that if  $X \succeq 0$  then  $\arcsin[X] - X \succeq 0$ .

[Hint: Use the fact that  $\arcsin(x) = \sum_{k=0}^{\infty} \frac{\binom{2k}{k}}{4^k(2k+1)} x^{2k+1}$  for  $x \in [-1, 1]$ ].

(d) Using the positive semidefinite assumption on  $A$  show then that  $\text{Tr}(A \arcsin[X]) \geq \text{Tr}(AX)$ . Conclude that  $v^* \geq \frac{2}{\pi} p_{SDP}^*$ .

#### 4 Lovász $\vartheta$ number

Let  $G = (V, E)$  be a graph. Recall the definition of Lovász theta number from Lecture 10:

$$\begin{aligned} & \underset{x \in \mathbb{R}^n, X \in \mathbf{S}^n}{\text{maximise}} && \sum_{i=1}^n x_i \\ & \text{subject to} && X_{ii} = x_i \quad i \in V \\ & && X_{ij} = 0 \quad ij \in E \\ & && \begin{bmatrix} 1 & x^T \\ x & X \end{bmatrix} \succeq 0 \end{aligned} \tag{3}$$

(a) Show that the dual of (3) can be expressed as

$$\begin{aligned} & \text{minimise} && Z_{00} \\ & \text{subject to} && z_i = (1 + Z_{ii})/2 \quad \forall i \in V \\ & && Z_{ij} = 0 \quad \forall \{i, j\} \in \bar{E} \\ & && \begin{bmatrix} Z_{00} & z^T \\ z & Z \end{bmatrix} \succeq 0 \end{aligned} \tag{4}$$

where  $\bar{E} = \{\{i, j\} : i \neq j \text{ and } \{i, j\} \notin E\}$  is the complement of  $E$  [Hint: you can use the fact (without proof) that  $\begin{bmatrix} A & B^T \\ B & C \end{bmatrix} \succeq 0 \iff \begin{bmatrix} A & -B^T \\ -B & C \end{bmatrix} \succeq 0$ ].

(b) Show that (4) can be simplified to:

$$\begin{aligned} & \text{minimise} && Z_{00} \\ & \text{subject to} && Z_{ii} = 1 \quad \forall i \in V \\ & && Z_{ij} = 0 \quad ij \in \bar{E} \\ & && \begin{bmatrix} Z_{00} & \mathbf{1}^T \\ \mathbf{1} & Z \end{bmatrix} \succeq 0 \end{aligned} \tag{5}$$

where  $\mathbf{1}$  denotes the vector with all ones [Hint: given  $(z, Z)$  feasible for (4), consider  $\tilde{Z}_{ij} = Z_{ij}/(z_i z_j)$ ].

(c) Use Slater condition to verify that (5) and (3) have the same optimal values.

(d) Show that for any graph  $G$  with  $n$  vertices we have  $\vartheta(G)\vartheta(\bar{G}) \geq n$  where  $\bar{G} = (V, \bar{E})$ .

## 5 Faces of the positive semidefinite cone

- (a) Recall the definition of a *face* of a convex set. Let  $V$  be a subspace of  $\mathbb{R}^n$ . Show that

$$F_V = \{Y \in \mathbf{S}_+^n : \text{im } Y \subseteq V\}$$

is a face of the positive semidefinite cone. What is its dimension?

- (b) Find a  $C \in \mathbf{S}^n$  such that  $\text{argmin}_{X \in \mathbf{S}_+^n} \langle C, X \rangle = F_V$  (this shows that  $F_V$  is an *exposed face* of  $\mathbf{S}_+^n$ ).
- (c) Let  $X \in \mathbf{S}_+^n$ . Show that the smallest face of  $\mathbf{S}_+^n$  containing  $X$  is  $F_{\text{im } X}$ .

## 6 Existence of extreme points

Given a set  $C \subseteq \mathbb{R}^n$  we say that  $C$  contains a *straight line* if there exists  $x \in C$  and  $v \in \mathbb{R}^n$  such that  $x + tv \in C$  for all  $t \in \mathbb{R}$ .

- (a) Let  $C$  be a nonempty closed convex set that does not contain any straight lines. Show that  $C$  has an extreme point [*Hint: you can use an argument by induction on the dimension of  $C$ , similar to the proof of Theorem 1.2 we did in lecture*].
- (b) Conversely, show that if  $C$  is a closed convex set with an extreme point then it does not contain any straight lines.

## 7 Extreme points in linear programming

- (a) Recall the definition of *extreme point* of a convex set.
- (b) Let  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$  and consider the convex set  $P = \{x \in \mathbb{R}_+^n : Ax = b\}$ . Show that any extreme point  $x$  of  $P$  satisfies  $|\text{supp}(x)| \leq m$  where  $\text{supp}(x) := \{i \in [n] : x_i \neq 0\}$  [*Hint: Show that if  $x$  is an extreme point of  $P$  then  $\ker(A) \cap \{y \in \mathbb{R}^n : \text{supp}(y) \subseteq \text{supp}(x)\} = \{0\}$ ].*]  
Use Exercise 6 to show that if  $P$  is not empty then it has at least one extreme point.
- (c) Use the result of part (b) to prove Carathéodory's theorem:

*Carathéodory's theorem:* Let  $S \subset \mathbb{R}^N$  be a finite set. Then any element of  $\text{conv}(S)$  can be expressed as a convex combination of at most  $N + 1$  points of  $S$ .

## 8 Extreme points in semidefinite programming

Part (a) of this exercise is the analogue of Exercise 7(a) for the case of semidefinite programming.

- (a) Let  $\mathcal{A} : \mathbf{S}^n \rightarrow \mathbb{R}^m$  be a linear map,  $b \in \mathbb{R}^m$  and let  $C = \{X \in \mathbf{S}_+^n : \mathcal{A}(X) = b\}$ . Show that any extreme point  $X$  of  $C$  satisfies  $r(r + 1)/2 \leq m$  where  $r = \text{rank } X$  [*Hint: Show that if  $X$  is an extreme point of  $C$  then  $\ker(\mathcal{A}) \cap \{Y \in \mathbf{S}^n : \text{im}(Y) \subseteq \text{im}(X)\} = \{0\}$ ].*]  
Use Exercise 6 to show that if  $C$  is nonempty then it has at least one extreme point.

- (b) Let  $A, B \in \mathbf{S}^n$ . Use part (a) to show that the set

$$R(A, B) = \{(x^T A x, x^T B x) : x \in \mathbb{R}^n\} \subseteq \mathbb{R}^2$$

is convex. (This set is known as the *numerical range* or *field of values* of the pair  $(A, B)$ .)  
 [Hint: consider  $\{(\langle A, X \rangle, \langle B, X \rangle) : X \in \mathbf{S}_+^n\}$ ].

- (c) Prove the following result, known as the *S-lemma*: Let  $A, B \in \mathbf{S}^n$  and assume that for any  $x \in \mathbb{R}^n$ ,  $x^T A x \geq 0 \Rightarrow x^T B x \geq 0$ . Assume furthermore that there exists  $z \in \mathbb{R}^n$  such that  $z^T A z > 0$ . Show that there exists  $\lambda \geq 0$  such that  $B \succeq \lambda A$ .

Give an example of  $A, B \in \mathbf{S}^2$  to show that the condition of existence of  $z \in \mathbb{R}^n$  such that  $z^T A z > 0$  cannot be removed in general.

## 9 Matrix square root

- (a) Let  $A, B \succ 0$ . Show that if  $A^2 \succeq B^2$  then  $A \succeq B$  [Hint: let  $v$  be an eigenvector of  $A - B$  and consider  $v^T (A + B)(A - B)v$ ].
- (b) Give an example of  $A, B \in \mathbf{S}_{++}^2$  such that  $A \succeq B$  but  $A^2 \not\succeq B^2$ .

## 10 Some facts about nonnegative polynomials

- (a) Show that if  $p \in \mathbb{R}[x_1, \dots, x_n]$  is nonnegative then it has even degree.
- (b) Show that if  $p = \sum_k q_k^2$  then necessarily  $\deg q_k \leq (\deg p)/2$ .

## 11 Homogeneous and nonhomogeneous polynomials

A polynomial  $p \in \mathbb{R}[x_1, \dots, x_n]$  is called *homogeneous of degree  $d$*  if it only involves monomials of degree exactly  $d$ . Given a nonhomogeneous polynomial  $p$  of degree  $d$  we can *homogenise* it by introducing an additional variable  $x_0$  via

$$\bar{p}(x_0, x_1, \dots, x_n) = x_0^d p(x_1/x_0, \dots, x_n/x_0) \quad (6)$$

- (a) Show that (6) is well-defined. What is the homogenisation of  $p(x_1, x_2) = x_1^2 x_2^2 - 2x_1 x_2 + 1$ ?
- (b) Show that  $p$  is nonnegative if and only if  $\bar{p}$  is nonnegative.
- (c) Show that  $p$  is a sum of squares if and only if  $\bar{p}$  is a sum of squares.

## 12 A nonnegative polynomial that is not a sum of squares

In lecture we saw the Motzkin polynomial  $M(x, y) = x^4 y^2 + x^2 y^4 + 1 - 3x^2 y^2$  which is an explicit example of a nonnegative polynomial that is not a sum of squares in the case  $(n, 2d) = (2, 6)$  (where  $n$  is the number of variables and  $2d$  the degree). In this exercise we look at a polynomial in 3 variables of degree 4 (i.e.,  $(n, 2d) = (3, 4)$ ) that is nonnegative but not a sum-of-squares. Consider the following polynomial (due to Choi and Lam [CL77]).

$$Q(x, y, z) = x^2 y^2 + x^2 z^2 + y^2 z^2 + 1 - 4xyz.$$

- (a) Show that  $Q(x, y, z) \geq 0$  for all  $(x, y, z) \in \mathbb{R}^3$ .
- (b) Show that  $Q$  is not a sum of squares.

### 13 Positive and decomposable maps

(Based on exercise 3.178 in [BPT12]) A map  $\Lambda : \mathbf{S}^{n_1} \rightarrow \mathbf{S}^{n_2}$  is called *positive* if  $\Lambda(A) \succeq 0$  whenever  $A \succeq 0$ .

- (a) Show that if  $\Lambda$  has the form  $\Lambda(A) = \sum_{i=1}^r P_i^T A P_i$  where  $P_1, \dots, P_r \in \mathbb{R}^{n_1 \times n_2}$  then  $\Lambda$  is positive. Such maps are called *decomposable*.
- (b) To any linear map  $\Lambda : \mathbf{S}^{n_1} \rightarrow \mathbf{S}^{n_2}$  we can consider the polynomial  $p(x, y) = y^T \Lambda(xx^T)y$  where  $x \in \mathbb{R}^{n_1}$  and  $y \in \mathbb{R}^{n_2}$ . Show that  $\Lambda$  is a positive map if and only if  $p$  is nonnegative. Show that  $\Lambda$  is decomposable if and only if  $p$  is a sum-of-squares.
- (c) Consider the following map  $\Lambda : \mathbf{S}^3 \rightarrow \mathbf{S}^3$  due to M.-D. Choi [Cho75]:

$$\Lambda(A) = 2 \begin{bmatrix} a_{11} + a_{22} & 0 & 0 \\ 0 & a_{22} + a_{33} & 0 \\ 0 & 0 & a_{33} + a_{11} \end{bmatrix} - A.$$

- (i) Show that  $\Lambda$  is positive [Hint: in the case  $a_{33} \geq a_{11}$  use  $\Lambda(A) = DAD + \begin{bmatrix} 2a_{22} & -2a_{12} & 0 \\ -2a_{12} & 2a_{33} & 0 \\ 0 & 0 & 2a_{11} \end{bmatrix}$  with  $D = \text{diag}(1, 1, -1)$ ; then generalise using cyclic symmetry of  $\Lambda$ ].
- (ii) Show that  $\Lambda$  is not decomposable. [Hint: show that the associated polynomial  $p(x, y)$  is not a sum-of-squares].

### 14 Sum-of-squares on the hypercube

(Based on [Ble15]) Let  $s(x) = x_1 + \dots + x_n$ .

- (a) Show that the function  $f(x) = (n - s(x))(n - 2 - s(x))$  is nonnegative on  $\{-1, 1\}^n$ .
- (b) Show that  $f$  is not 1-sos on  $\{-1, 1\}^n$ .
- (c) Show that  $f$  is 2-sos on  $\{-1, 1\}^n$  [Hint: what is  $(1 - x_i - x_j + x_i x_j)^2$  ?]

### References

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