## 10 Nonnegative polynomials, sums of squares and semidefinite programming

**Definition 10.1.** We say a polynomial  $p \in \mathbb{R}[x]$  is nonnegative if  $p(x) \geq 0$  for all  $x \in \mathbb{R}$ .

Note that if we have a "good" way of checking nonnegativity of polynomials then we can also minimise (or maximise) polynomials. Indeed if p(x) is a polynomial then

$$\min_{x \in \mathbb{R}} p(x) = \max_{x \in \mathbb{R}} \gamma \quad \text{s.t.} \quad p - \gamma \text{ is nonnegative.}$$
 (1)

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It is not difficult to verify that if p is nonnegative then:

- The degree of p is even and the leading coefficient (i.e., the coefficient of  $x^{2d}$  if deg p=2d) is nonnegative.
- Any real root of p has even multiplicity.

These conditions can also be shown to be sufficient (see proof of Theorem 10.1 below).

**Definition 10.2.** We say that a polynomial  $p \in \mathbb{R}[x]$  is a sum-of-squares if there exist polynomials  $q_1, \ldots, q_k \in \mathbb{R}[x]$  such that  $p = \sum_{i=1}^k q_i^2$ .

It is clear that if p is a sum of squares, then it is nonnegative. The converse is also true for polynomials in one variable.

**Theorem 10.1.** A univariate polynomial  $p(x) = \sum_{k=0}^{2d} p_k x^k$  of degree 2d is nonnegative if and only if there exist  $q_1, q_2$  of degree  $\leq d$  such that  $p = q_1^2 + q_2^2$ .

*Proof.* The implication  $\Leftarrow$  is clear. Assume p(x) is nonnegative. Since p has real coefficients, if p(z) = 0 then  $p(\bar{z}) = 0$ . Furthermore if z is a real root of p then it must have even multiplicity. This implies that we can write:

$$p(x) = p_{2d} \prod_{i=1}^{d} (x - z_i)(x - \bar{z_i}) = |q(x)|^2$$

where  $q(x) = \sqrt{p_{2d}} \prod_{i=1}^{d} (x-z_i)$  (note that  $p_{2d} \ge 0$  since p is nonnegative). If we let  $q_1(x) = \text{Re}[q(x)]$  and  $q_2 = \text{Im}[q(x)]$  (one can easily verify that these are polynomials of degree at most d) we get the desired result.

The next theorem shows that checking if a polynomial is a sum of squares can be decided using a semidefinite feasibility program:

**Theorem 10.2.** A polynomial  $p(x) = \sum_{k=0}^{2d} p_k x^k$  is a sum of squares if, and only if, there exists a positive semidefinite matrix M of size  $(d+1) \times (d+1)$  such that

$$p_k = \sum_{\substack{0 \le i, j \le d \\ i+j=k}} M_{ij} \qquad \forall k = 0, \dots, 2d.$$
(2)

(The rows and columns of the matrix M are indexed by  $0, \ldots, d$  instead of  $1, \ldots, d+1$  for convenience.)

*Proof.* We first prove  $\Rightarrow$ : assume p is a sum of squares, i.e.,  $p = \sum_{i=1}^{\ell} q_i^2$  where  $\deg q_i \leq d$ . Denote by  $[q_i] \in \mathbb{R}^{d+1}$  the vector of coefficients of  $q_i$  and  $[x] = (1, x, \dots, x^d)$  so that  $q_i(x) = [q_i]^T[x]$ . Then we have

$$p(x) = \sum_{i=1}^{\ell} q_i(x)^2 = \sum_{i=1}^{\ell} [x]^T [q_i] [q_i]^T [x] = [x]^T M[x]$$

where  $M = \sum_{i=1}^{\ell} [q_i][q_i]^T$ . Note that  $M \succeq 0$ . To see that the linear equations (2) hold, note that since  $p(x) = [x]^T M[x]$  and

$$[x]^{T}M[x] = \sum_{0 \le i, j \le d} M_{ij}x^{i}x^{j} = \sum_{k=0}^{2d} \left(\sum_{\substack{0 \le i, j \le d\\i+j=k}} M_{ij}\right)x^{k}$$
(3)

we must have (by matching coefficients)  $p_k = \sum_{\substack{0 \le i,j \le d \\ i+j=k}} M_{ij}$  for all  $k=0,\ldots,2d$ . We now prove the converse. Assume  $M \succeq 0$  satisfies (2). Then we can decompose M as  $M = \sum_{i=1}^{\ell} [q_i][q_i]^T$  for some vectors  $[q_i] \in \mathbb{R}^{d+1}$ . Define  $q_i(x) = [q_i]^T[x]$  where  $[x] = (1, x, \ldots, x^d)$ . Then one can easily verify that  $\sum_{i=1}^{\ell} q_i(x)^2 = [x]^T M[x] = p(x)$  where the last equality follows from (2) and by the same calculation as in (3).

By combining Theorems 10.1 and 10.2 we can check if a polynomial in one variable is nonnegative using a semidefinite feasibility problem. In fact using (1) we can minimize any polynomial in one variable using semidefinite programming:

 $\min_{x \in \mathbb{R}} p(x) = \max \quad \gamma \quad \text{s.t.} \quad p - \gamma \text{ nonnegative}$ 

$$= \max_{\gamma \in \mathbb{R}, M \in \mathbf{S}^{d+1}} \quad \gamma \quad \text{s.t.} \quad \begin{cases} M \succeq 0 \\ p_k = \sum_{\substack{0 \le i, j \le d \\ i+j=k}} M_{ij} & \forall k = 1, \dots, 2d \\ p_0 - \gamma = M_{00} \end{cases}$$

**Example 1** (Polynomials of degree 2). We know from high-school algebra that a polynomial p(x) = 0 $ax^2 + bx + c$  is nonnegative iff  $b^2 - 4ac \le 0$  and  $a, c \ge 0$ . Theorems 10.1 and 10.2 tell us that this polynomial is nonnegative if and only if there exists a matrix  $M \in \mathbf{S}^2$  such that

$$M \succeq 0,$$
  
 $M_{00} = c,$   
 $M_{01} + M_{10} = b,$   
 $M_{11} = a.$ 

This is equivalent to saying that  $\begin{bmatrix}c&b/2\\b/2&a\end{bmatrix}$  is positive semidefinite, which in turn is equivalent to having  $b^2 - 4ac \le 0$  and  $a, c \ge 0$ .

Let  $P_{2d}$  be the cone of nonnegative polynomials (in one variable) of degree 2d:

$$P_{2d} = \left\{ (p_0, \dots, p_{2d}) \in \mathbb{R}^{2d+1} : \sum_{k=0}^{2d} p_k x^k \ge 0 \ \forall x \in \mathbb{R} \right\}.$$

**Theorem 10.3.**  $P_{2d}$  is a proper cone.

*Proof.* We have to check that  $P_{2d}$  is closed, convex, pointed and has nonempty interior. Note that  $P_{2d}$  can be written as

$$P_{2d} = \bigcap_{x \in \mathbb{R}} \underbrace{\left\{ (p_0, \dots, p_{2d}) : \sum_{k=0}^{2d} p_k x^k \ge 0 \right\}}_{H_x}$$

where each  $H_x$  is a closed halfspace. Thus  $P_{2d}$  is closed and convex as an intersection of closed convex sets. Checking that  $P_{2d}$  is pointed is easy. We leave it as an exercise to verify that the polynomial  $x^{2d} + 1$  is in the interior of  $P_{2d}$ .