Equivariant semidefinite lifts and sum-of-squares hierarchies

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ISMP 2015

Central question in optimization is to optimize a linear function ℓ on a convex set C:

 $\min_{x\in C}\ell(x).$

- Need "good" description of C to solve optimization problem efficiently.
- Conic programming descriptions

Semidefinite programming representations

Semidefinite programming

min
$$\mathcal{L}(Y)$$
 subject to $Y \in \mathbf{S}^d_+, Y \in L$

where L affine subspace of S^d .

Positive semidefinite lift of *C*: $C = \pi(\mathbf{S}^d_+ \cap L)$

where

- π linear map
- L affine subspace of S^d



Consequence: Optimizing linear function over *C* is SDP of size *d*.

$$\min_{x \in C} \ell(x) = \min (\ell \circ \pi)(Y) \text{ s.t. } Y \in \mathbf{S}^d_+ \cap L$$

Example of psd lift

$$[-1,1]^{2} = \left\{ (x_{1},x_{2}) \in \mathbb{R}^{2} : \exists u \in \mathbb{R} \ \begin{bmatrix} 1 & x_{1} & x_{2} \\ x_{1} & 1 & u \\ x_{2} & u & 1 \end{bmatrix} \succeq 0 \right\}$$
(1)



Sum-of-squares lifts

- Let $X \subseteq \mathbb{R}^n$. Goal: find a psd lift of $P = \operatorname{conv}(X)$.
- Let $\mathcal{F}(X, \mathbb{R})$ be the space of real-valued functions on *X*.

Theorem (Lasserre, Gouveia et al.)

Assume there is a subspace V of $\mathcal{F}(X, \mathbb{R})$ such that any facet-defining inequality $\ell(x) \leq b$ for $\operatorname{conv}(X)$ can be certified using sum-of-squares in V, i.e., there exist $f_1, \ldots, f_J \in V$:

$$b-\ell(x)=\sum_{j=1}^J f_j(x)^2 \quad orall x\in X.$$

Then conv(X) has an (explicit) psd lift of size dim V.



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 Lasserre / theta-body hierarchy: take V = Pol_{≤k}(X), subspace of polynomials of degree at most k on X.

Psd lifts that respect symmetry

• Assume that $X \subset \mathbb{R}^n$ is invariant under action of group *G*.

We are interested in finding a psd lift of P = conv(X) that **respects the** *G*-symmetry of *P*.

- Example of polytopes with symmetries:
 - Regular *N*-gon in \mathbb{R}^2
 - Hypercube [-1,1]ⁿ
 - Combinatorial polytopes: cut polytope, matching polytope, etc.



Equivariant psd lifts: definition

PSD lift

$$C = \pi(\mathbf{S}^d_+ \cap L)$$

respects *G*-symmetry of *C* if any transformation $g \in G$ can be lifted to a transformation $\Phi(g)$ upstairs such that

$$\pi(\Phi(g)Y) = g\pi(Y)$$

for all $g \in G$, $Y \in \mathbf{S}^d_+ \cap L$.



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for all $g \in G, Y \in \mathbf{S}^d_+ \cap L$.

Requirements on Φ :

• $\Phi(g)$ acts on **S**^d by congruence transformation, i.e.,

$$\Phi(g) Y = \rho(g) Y \rho(g)^{T}$$

where $\rho: G \to GL_d(\mathbb{R})$ homomorphism

• $\Phi(g)$ leaves *L* invariant

(These two conditions imply that $\Phi(g)$ leaves $\mathbf{S}^d_+ \cap L$ invariant)

A psd lift that respects G-symmetry of P is called equivariant psd lift





Sum-of-squares lifts that are equivariant

• Action of *G* on *X* induces action on $\mathcal{F}(X, \mathbb{R})$:

 $(g \cdot f)(x) = f(g^{-1}x)$

• Subspace $V \subset \mathcal{F}(X, \mathbb{R})$ is called *G*-invariant if $G \cdot V = V$.

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Theorem (Sum-of-squares lifts that are equivariant)

Assume there is a *G*-invariant subspace *V* of $\mathcal{F}(X, \mathbb{R})$ such that any facet-defining inequality $\ell(x) \leq b$ for conv(*X*) can be certified using sum-of-squares in *V*, *i.e.*,

$$b-\ell(x)=\sum_{j=1}^J f_j(x)^2$$
 where $f_j\in V$ $(j=1,\ldots,J).$

Then conv(X) has an (explicit) G-equivariant psd lift of size dim V.

• Lasserre/theta-body lifts are equivariant because $V = Pol_{\leq k}(X)$ is *G*-invariant.

Main result: structure theorem

• Let $G \subset GL_n(\mathbb{R})$ and $X = G \cdot x_0$ where $x_0 \in \mathbb{R}^n$, and

 $P = \operatorname{conv}(G \cdot x_0).$

Theorem (Structure theorem)

Assume that P has a G-equivariant psd lift of size d. Then there exists an G-invariant subspace V of $\mathcal{F}(X, \mathbb{R})$ with dim $V \leq d^3$, such that any facet-defining inequality $\ell(x) \leq b$ for P has a sum-of-squares certificate with functions from V, i.e.,

$$b-\ell(x)=\sum_{j=1}^J f_j(x)^2$$
 where $f_j\in V$ $(j=1,\ldots,J).$

Consequence of structure theorem: To study *G*-equivariant psd lifts of conv(*X*), need to study *G*-invariant subspaces of $\mathcal{F}(X, \mathbb{R})$.

In particular we will be interested in decomposing $\mathcal{F}(X, \mathbb{R})$ into a direct sum of irreducibles:

$$\mathcal{F}(X,\mathbb{R})=\bigoplus_i V_i$$

where each V_i is irreducible.

(Recall that subspace V is irreducible if it does not contain a nontrivial invariant subspace).

Example 1: regular polygons

$$X = \left\{ \left(\cos(2\pi k/N), \sin(2\pi k/N) \right) : k = 1, \dots, N \right\}$$



• $\mathcal{F}(X,\mathbb{R})$ decomposes as

$$\mathcal{F}(X,\mathbb{R}) = \mathsf{TPol}_0 \oplus \mathsf{TPol}_1 \oplus \cdots \oplus \mathsf{TPol}_{|N/2|}$$

where $\text{TPol}_k = \text{span}(c_k, s_k)$

$$\begin{cases} C_k(x_1, x_2) = \operatorname{Re}[(x_1 + ix_2)^k] \\ S_k(x_1, x_2) = \operatorname{Im}[(x_1 + ix_2)^k] \end{cases}$$

("Discrete Fourier Transform" for signals of length N)

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Fact

Any G-invariant subspace V of $\mathcal{F}(X, \mathbb{R})$ has the form $V = \bigoplus_{k \in K} \text{TPol}_k$ where $K \subseteq \{0, \dots, \lfloor N/2 \rfloor\}$.

Example 1: regular polygons (continued)

$$\mathcal{F}(X,\mathbb{R}) = \mathsf{TPol}_0 \oplus \mathsf{TPol}_1 \oplus \cdots \oplus \mathsf{TPol}_{\lfloor N/2 \rfloor}$$

Theorem

Assume $K \subseteq \{0, \ldots, \lfloor N/2 \rfloor\}$ is such that we can write:

$$\cos(\pi/N) - x_1 = \sum_{j=1}^J f_j(x)^2 \quad \forall x \in X.$$

where $f_1, \ldots, f_J \in \bigoplus_{k \in \mathbf{K}} \mathsf{TPol}_k$. Then $|\mathbf{K}| \ge \ln(N/2)/2$.

Consequence: any equivariant psd lift of the regular *N*-gon must have size at least $\Omega(\log N)$. Bound is tight (cf. James's talk tomorrow).

(2)

Example 2: parity polytope

Parity polytope:

$$PAR_n = conv(EVEN_n)$$

where

$$EVEN_n = \left\{ x \in \{-1, +1\}^n : \prod_{i=1}^n x_i = 1 \right\}.$$

- PAR_n is symmetric with respect to:
 - Switching the sign of an even number of components (*G*_{even} symmetry)
 Permutation of components (*G*_n symmetry).

Symmetry group of PAR_n

$$\Gamma = evenly signed permutations$$

$$= G_{\text{even}} \rtimes \mathfrak{S}_n.$$

Example 2: parity polytope (continued)

Lemma (1)

• $\mathcal{F}(EVEN_n, \mathbb{R})$ decomposes into Γ -irreducibles as:

$$\mathcal{F}(EVEN_n, \mathbb{R}) = \operatorname{Pol}_0 \oplus \operatorname{Pol}_1 \oplus \cdots \oplus \operatorname{Pol}_{\lfloor n/2 \rfloor}$$
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where Pol_k is the space of square-free polynomials of degree at most k.

$$\operatorname{dim} \operatorname{Pol}_{k} = \begin{cases} \binom{n}{k} & \text{if } k < n/2\\ \frac{1}{2} \binom{n}{n/2} & \text{if } k = n/2. \end{cases}$$

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The Lasserre/theta-body hierarchy for PAR_n requires at least $\lceil n/4 \rceil$ steps.

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Corollary (Structure theorem + Lemma 1 + Lemma 2)

Any Γ -equivariant psd lift for PAR_n has size at least $\binom{n}{\lceil n/4 \rceil}$.

Conclusion

 Main message: to study equivariant psd lifts, need to study invariant subspaces of *F*(*X*, ℝ).

Lower bounds

- $\Omega(\log N)$ lower bound on equivariant psd lifts of regular *N*-gon.
- Exponential lower bound on equivariant psd lifts of parity polytope, *via* lower bound on Lasserre/theta-body hierarchy. Similar analysis also gives exponential lower bound for equivariant psd lifts of cut polytope.
- Upper bounds: Understanding invariant subspaces of *F*(*X*, ℝ) is also useful for constructions/upper bounds → sparse sum-of-squares certificates (see James's talk tomorrow FB04).

Papers: arXiv:1312.6662 (parity polytope + cut polytope) arXiv:1409.4379 (regular polygons)

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