# Polyhedral approximations of the positive semidefinite cone

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# Conic programming

min. 
$$\langle c, x \rangle$$
  
s.t.  $A(x) = b, x \in K$ 

K is a convex cone, and A linear map

- Three major cones
  - Linear programming (LP):  $K = \mathbb{R}^m_+$
  - Second-order cone programming (SOCP):  $K = Q^{m_1} \times \cdots \times Q^{m_k}$  where

$$\mathcal{Q}^m = \{(x,t) \in \mathbb{R}^m \times \mathbb{R} : ||x||_2 \leq t\}.$$

 Semidefinite programming (SDP): K = S<sup>m</sup><sub>+</sub> (m × m symmetric positive semidefinite matrices)

### Reduction between conic programs

It is known that

$$``LP \subsetneq SOCP \subsetneq SDP''$$

- $LP \neq SOCP, SDP$  because  $\mathcal{Q}^m$  and  $\mathbf{S}^m_+$  are nonpolyhedral
- Not possible to express  $S^3_+$  using second-order cones! [F17]

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#### What about approximations?

• Ben-Tal & Nemirovski:  $Q^m$  can be  $\epsilon$ -approximated using an LP of size  $O(1)m\log(1/\epsilon)$ . More precisely:

#### Theorem (Ben-Tal & Nemirovski, 2001)

For any  $\epsilon > 0$ , there is a polytope P with extension complexity  $\leq O(1)m\log(1/\epsilon)$  such that  $(1-\epsilon)B^m \subset P \subset B^m$ , where  $B^m$  is the unit Euclidean ball in  $\mathbb{R}^m$ 

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• What about the PSD cone?

$$\begin{split} \mathcal{C} &= \text{convex set of } n \times n \text{ real symmetric} \\ & \text{positive semidefinite matrices of trace 1} \\ &= \text{conv} \left\{ xx^{\mathcal{T}} : x \in \mathbb{R}^n, \|x\|_2 = 1 \right\}. \end{split}$$

**Question:** How well can we approximate *C* using polyhedra?



# Extension complexity of polytopes

A polytope P has extension complexity N if it can be written as

 $P = \pi(Q) \quad \text{where} \quad \begin{cases} Q = \{x \in \mathbb{R}^{N} : x \ge 0 \text{ and } Ax = b\} \\ \pi \text{ linear (projection) map} \end{cases}$ 



Note:

- $xc(P) \le \#vertices(P)$  and  $xc(P) \le \#facets(P)$ .
- Also  $xc(P) = xc(P^*)$  (invariant under duality)

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• Braun, Fiorini, Pokutta, Steurer 2013: There exists a spectrahedral shadow S in  $\mathbb{R}^{n^2}$  such that rank<sub>psd</sub> $(S) \leq n + 1$  and whenever  $S \subset P \subset S + \epsilon B_1$  then  $\operatorname{xc}(P) \geq e^{cn}$ , where  $B_1$  is the  $\ell_1$  ball.

# Gaussian width

Given a direction u, the *width* of a set S in this direction is:

$$w(S, u) = \max_{x \in S} \langle u, x \rangle - \min_{x \in S} \langle u, x \rangle.$$

The Gaussian width of S is

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#### Theorem

Assume P is a polytope such that  $C \subset P$  and  $w_G(P) \leq 2w_G(C)$ . Then  $xc(P) \geq e^{cn^{1/3}}$ .

## Extension complexity: Yannakakis theorem

Controlling extension complexity is much more difficult than # vertices. Very useful tool is Yannakakis theorem...

Let  $C \subset D$  be two nested convex sets.



• Slack matrix of D wrt C is an (infinite) matrix S

$$S[x,\ell] = 1 - \ell(x)$$

where

- x = extreme point of *C*
- $\ell = \text{extreme point of } D^{\circ} \text{ (polar)}$

#### Theorem (Yannakakis 1991)

Let  $C \subset D$  compact convex sets. There exists a polytope P with xc(P) = N s.t.  $C \subset P \subset D$  iff the slack matrix S has a nonnegative factorization of size N.

Nonnegative factorization:

$$S[x,\ell] = \sum_{i=1}^{N} a_i(x)b_i(\ell)$$

where  $a_i(x), b_i(\ell) \geq 0$ .

#### Our slack matrix

$$C = \operatorname{conv}\left\{xx^{T} : \|x\|_{2} = 1\right\}$$

• Can easily show that slack matrix of C wrt  $(1-\epsilon)C$  is

$$S[x,y] = (1-\epsilon)n(x^Ty)^2 + \epsilon \quad \forall x,y \in S^{n-1}.$$

• A nonnegative factorization of S of size N:

$$S[x,y] = \sum_{i=1}^N f_i(x)g_i(y) \quad orall x,y \in S^{n-1}.$$

where  $f_i, g_i \ge 0$ 

• Useful normalization:  $\int_{S^{n-1}} f_i = 1$  and  $\sum_{i=1}^{N} g_i \equiv 1$ . (S is a "column-stochastic" matrix)

#### Interpretation

$$\underbrace{(1-\epsilon)n(x^{T}y)^{2}+\epsilon}_{\tilde{Q}_{y}(x)} = \sum_{i=1}^{N} f_{i}(x)g_{i}(y) \quad \forall x, y \in S^{n-1}$$

•  $ilde Q_y$  is a quadratic form with  $ilde Q_y(x) \geq \epsilon$  for  $x \in S^{n-1}$  and  $\int ilde Q_y(x) dx = 1$ 

- Interpretation of nonnegative factorization: All Q<sub>y</sub>'s are convex combinations of the functions {f<sub>1</sub>,..., f<sub>N</sub>}.
- To prove a lower bound, we need to show that we need many functions  $f_i$ 's to cover all the  $\tilde{Q}_y$ 's  $(y \in S^{n-1})$
- Important: the functions  $f_i$  need not be quadratic!



Assume  $f_i$  are quadratic s.t.  $\tilde{Q}_y \in \text{conv}(f_1, \ldots, f_N)$  for all  $y \in S^{n-1}$ . Then  $N \geq \exp(cn)$ .



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- Evaluate at x = y:  $\tilde{Q}_y(y) = (1 \epsilon)n + \epsilon \in n \operatorname{conv}((a_1^T y)^2, \dots, (a_N^T y)^2).$



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- Evaluate at x = y:  $\tilde{Q}_y(y) = (1 \epsilon)n + \epsilon \in n \operatorname{conv}((a_1^T y)^2, \dots, (a_N^T y)^2).$
- Implies: for any  $y \in S^{n-1}$ , there exists  $i \in \{1, ..., N\}$  s.t.  $(a_i^T y)^2 \ge 1 \epsilon = 2/3$
- Hence  $N \ge \exp(cn)$

# Discretizing the sphere $S^{n-1}$

Spherical cap around  $a \in S^{n-1}$ :

 $\{x \in S^{n-1} : \langle a, x \rangle \ge 2/3\}.$ 

Surface area of this spherical cap is  $\leq e^{-cn}!$ 

Need exponentially many such spherical caps to cover sphere!



• Fourier decomposition on  $S^{n-1}$ : any function  $f: S^{n-1} \to \mathbb{R}$  has a Fourier decomposition

$$f = Y_0 + Y_1 + Y_2 + \dots = \sum_{k=0}^{\infty} Y_k$$

(similar to Fourier decomposition on the cube)

• The  $Y_k$  are harmonic polynomials of degree k.

### A "low-pass" filter

If  $f = \sum_{k=0}^{\infty} Y_k$  we can apply a "low-pass filter" (smoothing operation) to f:

$$P_{\rho}f = \sum_{k=0}^{\infty} \rho^k Y_k$$



where 0  $\leq \rho < 1$  (Poisson kernel).



We want to quantify how much  $P_{\rho}$  flattens functions.

### $L^p$ norms

• Given  $f:S^{n-1} \to \mathbb{R}$  and  $p \ge 1$  define

$$\|f\|_p = \int_{S^{n-1}} |f|^p d\sigma.$$

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• Hölder: 
$$||f||_1 \le ||f||_2 \le \cdots \le ||f||_{\infty}$$
.  
• Ratio  $\frac{||f||_q}{||f||_p}$  for  $q > p$  tells us how flat/spiked  $f$  is.

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Example: If f poly of degree d then  $||f||_p \leq (p-1)^{d/2} ||f||_2$  for  $p \geq 2$ 

### Properties of low-pass filter $P_{\rho}$

$$P_{\rho}f = \sum_{k=0}^{\infty} \rho^k Y_k$$

Contractivity:  $P_{\rho}f$  is no more spiked than f:

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Theorem (Hypercontractivity of Poisson kernel (Beckner))

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# Theorem (Hypercontractivity of Poisson kernel (Beckner)) For any $p \ge 1$ , $||P_{\rho}f||_q \le ||f||_p$ for $q = 1 + \rho^{-2}(p-1) > p$ .

- Hypercontractivity: Nelson 1960s for a smoothing kernel in Gaussian space.
- Gross 1975: logarithmic Sobolev inequalities
- Generalization to other smoothing kernels. Applications in computer science. [de Wolf, O'Donnell, Klartag-Regev].

# A lemma

#### Lemma

Let  $f: S^{n-1} \to \mathbb{R}$  s.t.  $f \ge 0$ ,  $\int f = 1$  and max  $f \le e^{\sqrt{n}}$ . Let  $\rho = \sqrt{5/n}$  (smoothing kernel parameter). Then

$$\sigma\left\{P_{\rho}f\geq4\right\}\leq c^{-\sqrt{n}}$$

for some absolute constant c > 1.

#### Proof.

Markov's inequality + hypercontractivity with good choices of q, p.

• Assume that  $f_1,\ldots,f_N:S^{n-1}\to\mathbb{R}_+$ ,  $\int f_i=1$  are such that:

$$\forall e \in S^{n-1}, \ (1-\epsilon)Q_e + \epsilon \in \operatorname{conv}(f_1, \ldots, f_N)$$

where  $Q_e(x) = n(e^T x)^2$  [extreme rays of the psd cone].

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• Apply  $P_{
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• Consequence: for any  $e \in S^{n-1}$  there is at least one *i* s.t.  $(P_{\rho}f_i)(e) \ge 4$ 

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- Previous lemma (+some technical details) tells us  $N \ge \exp(c\sqrt{n})$ .

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