Numerical analysis in the age of data science through the Solvability Complexity Index Hierarchy

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When can algorithms be trusted?

- Ex1: Data-driven dynamical systems Error control and verification?
- Ex2: Trustworthiness in AI Stability and accuracy guarantees?

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Why is this question essential?

- Bedrock of numerical analysis.
- Reliable computations to back-up and test theories.
- Computed "solutions" meaningless without understanding error.

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Toolkit for classifying the difficulty of computational problems and proving the optimality of algorithms.

Talk outline

- **Example 1:** Rigorous and verifiable data-driven dynamical systems
 - Upper bounds (on difficulty): Positive results on what's possible.
- A toolkit: Solvability Complexity Index Hierarchy
 - Classification for infinite-dimensional spectral problems and beyond.
- Example 2: Smale's 18th problem, "What are the limits of AI?"
 - Lower bounds (on difficulty): Use SCI techniques to prove impossibility results.

Toolkit for classifying the difficulty of computational problems and proving the optimality of algorithms.

Example 1: Rigorous and verifiable data-driven dynamical systems

Upper bounds: Positive results on what's possible.

Data-driven dynamical systems

• State $x \in \Omega \subseteq \mathbb{R}^d$, *unknown* function $F: \Omega \to \Omega$ governs dynamics

$$x_{n+1} = F(x_n)$$

- Goal: Learn about system from data $\{x^{(m)}, y^{(m)} = F(x^{(m)})\}_{m=1}^{M}$
 - E.g., data from trajectories, experimental measurements, simulations, ...
 - E.g., used for forecasting, control, design, understanding, ...
- Applications: chemistry, climatology, electronics, epidemiology, finance, fluids, molecular dynamics, neuroscience, plasmas, robotics, video processing, ...

Can we develop verified methods?



Operator viewpoint

• Koopman operator $\mathcal K$ acts on $\underline{\mathrm{functions}}\;g\colon\Omega\to\mathbb C$

 $[\mathcal{K}g](x) = g(F(x))$

• $\mathcal K$ is *linear* but acts on an *infinite-dimensional* space.



• Work in $L^2(\Omega, \omega)$ for positive measure ω , with inner product $\langle \cdot, \cdot \rangle$.

- Koopman, "Hamiltonian systems and transformation in Hilbert space," PNAS, 1931.
- Koopman, v. Neumann, "Dynamical systems of continuous spectra," PNAS, 1932.



Encodes: geometric features, invariant measures, transient behaviour, long-time behaviour, coherent structures, quasiperiodicity, etc.

GOAL: Data-driven approximation of $\mathcal K$ and its spectral properties.

[•] Mezić, "Spectral properties of dynamical systems, model reduction and decompositions," Nonlinear Dynamics, 2005.

Koopmania*: a revolution in the big data era

New Papers on "Koopman Operators"



 \approx 35,000 papers over last decade!

Very little on convergence guarantees or verification.

Why is this lacking?

- K. operators have been distinct from NA community.
- Dealing with infinite dim. is notoriously hard ...

*Wikipedia: "its wild surge in popularity is sometimes jokingly called 'Koopmania'"

Can we compute spectral properties in inf. dim.?

Spec(\mathcal{K}) = { $\lambda \in \mathbb{C}: \mathcal{K} - \lambda I$ is not invertible}

"Operators that arise in practice are not **diagonalized**, and it is often very hard to locate the spectrum. Thus, one has to settle for numerical approximations. Unfortunately, there are **no proven** <u>general</u> techniques." W. Arveson, Berkeley (1994)

Naïve: $\mathcal{K} \longrightarrow \mathbb{K} \in \mathbb{C}^{N \times N}$ + compute e-values, **problems:**

- **1)** "Too much": Approximate spurious modes $\lambda \notin \text{Spec}(\mathcal{K})$ "spectral pollution"
- **2)** "Too little": Miss parts of $Spec(\mathcal{K})$
- 3) Continuous spectra
- 4) Verification: Which part of an approximation can we trust?
- Arveson, "The role of C*-algebras in infinite dimensional numerical linear algebra," Contemp. Math., 1994.
- Davies, "Linear operators and their spectra," CUP, 2007.
- Brunton, Kutz, "Data-driven Science and Engineering: Machine learning, Dynamical systems, and Control," CUP, 2019.

Dynamic Mode Decomposition (DMD)

Given dictionary $\{\psi_1, \dots, \psi_N\}$ of functions $\psi_j \colon \Omega \to \mathbb{C}$

- Schmid, "Dynamic mode decomposition of numerical and experimental data," Journal of fluid mechanics, 2010.
- Kutz, Brunton, Brunton, Proctor, "Dynamic mode decomposition: data-driven modeling of complex systems," SIAM, 2016.
- Williams, Kevrekidis, Rowley "A data-driven approximation of the Koopman operator: Extending dynamic mode decomposition," Journal of Nonlinear Science, 2015.

 $\{x^{(m)}, y^{(m)} = F(x^{(m)})\}_{m}^{M}$

Residual DMD (ResDMD): Approx. \mathcal{K} and $\mathcal{K}^*\mathcal{K}$

$$\langle \psi_k, \psi_j \rangle \approx \sum_{m=1}^M w_m \overline{\psi_j(x^{(m)})} \psi_k(x^{(m)}) = \left[\underbrace{\Psi_X^* W \Psi_X}_{G} \right]_{jk}$$

$$\langle \mathcal{K}\psi_k, \psi_j \rangle \approx \sum_{m=1}^M w_m \overline{\psi_j(x^{(m)})} \underbrace{\psi_k(y^{(m)})}_{[\mathcal{K}\psi_k](x^{(m)})} = \left[\underbrace{\Psi_X^* W \Psi_Y}_{K_1} \right]_{jk}$$
Error control
$$\langle \mathcal{K}\psi_k, \mathcal{K}\psi_j \rangle \approx \sum_{m=1}^M w_m \overline{\psi_j(y^{(m)})} \psi_k(y^{(m)}) = \left[\underbrace{\Psi_Y^* W \Psi_Y}_{K_2} \right]_{jk}$$

$$\underbrace{ \text{Residuals:}}_{g} = \sum_{j=1}^N \mathbf{g}_j \psi_j, \ \|\mathcal{K}g - \lambda g\|^2 \approx \mathbf{g}^* [K_2 - \lambda K_1^* - \overline{\lambda} K_1 + |\lambda|^2 G] \mathbf{g}$$

- C., Townsend, "Rigorous data-driven computation of spectral properties of Koopman operators for dynamical systems,"
 Communications on Pure and Applied Mathematics, under review.
- C., Ayton, Szőke, "Residual Dynamic Mode Decomposition," Journal of Fluid Mechanics, under review.

Example of an upper bound

$$\operatorname{res}(\lambda, \mathbf{g}) = \sqrt{\frac{\mathbf{g}^*[K_2 - \lambda K_1^* - \overline{\lambda} K_1 + |\lambda|^2 G]\mathbf{g}}{\mathbf{g}^* G \mathbf{g}}}, \qquad \operatorname{Spec}_{\varepsilon}(\mathcal{K}) = \bigcup_{\|\mathcal{B}\| \le \varepsilon} \operatorname{Spec}(\mathcal{K} + \mathcal{B})$$
1. Compute $G, K_1, K_2 \in \mathbb{C}^{N \times N}$.
2. For z_k in comp. grid, compute $\tau_k = \min_{\substack{g = \sum_{j=1}^N \mathbf{g}_j \psi_j}} \operatorname{res}(z_k, g)$, corresponding g_k (gen. SVD).
3. Output: $\{z_k: \tau_k < \varepsilon\}$ (approx. of $\operatorname{Spec}_{\varepsilon}(\mathcal{K})$), $\{g_k: \tau_k < \varepsilon\}$ (ε -pseudo-eigenfunctions).

Theorem: Suppose the quadrature rule converges.

- Error control: $\{z_k: \tau_k < \varepsilon\} \subseteq \operatorname{Spec}_{\varepsilon}(\mathcal{K})$ (as $M \to \infty$)
- **Convergence:** Converges locally uniformly to $\operatorname{Spec}_{\varepsilon}(\mathcal{K})$ (as $N \to \infty$)

Similarly, overcome: 1) "too much", 2) "too little", 3) cts spectra, 4) verification. Paves the way for rigorous data-driven Koopmania!

Example: pressure field of turbulent flow



• Reynolds number $\approx 3.9 \times 10^5$

0.25

0.2

0.15

0.1

0.05

-0.05

-0.1

-0.15

-0.2

-0.25

0

 Ambient dimension ≈ 300,000 (number of measurement points)







Example: pressure field of turbulent flow



Previous algorithm Γ_M (with adaptive N = N(M)):

$$\lim_{\varepsilon \downarrow 0} \lim_{M \to \infty} \Gamma_M \left(\left\{ x^{(m)}, y^{(m)} = F(x^{(m)}) \right\}_{m=1}^M, \varepsilon \right) = \operatorname{Spec}(\mathcal{K})$$

Phenomena of "successive limits" widespread...

A toolkit: Solvability Complexity Index Hierarchy

Solvability Complexity Index Hierarchy

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For $A \in \Omega$, want to compute $\Xi: \Omega \to (\mathcal{M}, d)$ — metric space

- Δ_0 : Problems solved in finite time (v. rare).
- Δ_1 : Problems solved in "one limit" with full error control: $d(\Gamma_n(A), \Xi(A)) \le 2^{-n}$
- Δ_2 : Problems solved in "one limit" (SCI=1): $\lim_{n \to \infty} \Gamma_n(A) = \Xi(A)$
- Δ_3 : Problems solved in "two successive limits" (SCI=2):

$$\lim_{n\to\infty}\lim_{m\to\infty}\Gamma_{n,m}(A)=\Xi(A)$$

and so on...

- C., "The Foundations of Infinite-Dimensional Spectral Computations," Cambridge, PhD thesis.
- Ben-Artzi, C., Hansen, Nevanlinna, Seidel, "On the solvability complexity index hierarchy and towers of algorithms."
- Hansen, "On the solvability complexity index, the *n*-pseudospectrum and approximations of spectra of operators," JAMS, 2011.
- McMullen, "Families of rational maps and iterative root-finding algorithms," Annals of Mathematics, 1987.
- Smale, "The fundamental theorem of algebra and complexity theory," Bulletin of the AMS, 1981.

Error control for spectral problems



• $\Sigma_1 \subsetneq \Delta_2$: $\lim_{n \to \infty} \Gamma_n(A) = \Xi(A)$, $\max_{z \in \Gamma_n(A)} \operatorname{dist}(z, \Xi(A)) \le 2^{-n}$ • $\Pi_1 \subsetneq \Delta_2$: $\lim_{n \to \infty} \Gamma_n(A) = \Xi(A)$, $\max_{z \in \Xi(A)} \operatorname{dist}(z, \Gamma_n(A)) \le 2^{-n}$ NB: Typically, eigensolvers for PDEs prove Δ_2 , not optimal Σ_1 ...

Σ_1 example: e-values with guaranteed error bounds

$$A = -\nabla^2 + x^2 + V(x)$$
 on \mathbb{R}^1

V	$\cos(x)$	tanh(x)	$\exp(-x^2)$	$1/(1 + x^2)$
E ₀	1.7561051579	0.8703478514	1.6882809272	1.7468178026
E ₁	3.3447026910	2.9666370800	3.3395578680	3.4757613534
<i>E</i> ₂	5.0606547136	4.9825969775	5.2703748823	5.4115076464
<i>E</i> ₃	6.8649969390	6.9898951678	7.2225903394	7.3503220313
E_4	8.7353069954	8.9931317537	9.1953373991	9.3168983920

Can deal with mix of cts and discrete spectra, other domains, other PDEs etc.

Increasing difficulty



Increasing difficulty



Increasing difficulty



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Increasing difficulty



*Open problem of Schwinger: "The special canonical group," "Unitary operator bases," PNAS, 1960.

Increasing difficulty



*Open problem of Schwinger: "The special canonical group," "Unitary operator bases," PNAS, 1960.

Increasing difficulty



*Open problem of Schwinger: "The special canonical group," "Unitary operator bases," PNAS, 1960.

Some success stories in inf. dim. spec. comp.

• Spectra with Σ_1 error control (v. large class of disc. ops and PDEs)

- C., Roman, Hansen, "How to compute spectra with error control," PRL, 2019.
- C., Hansen, "The foundations of spectral computations via the solvability complexity index hierarchy," JEMS, under revisions.
- Spectral measures and type (classification varies from Δ_1 to Σ_3)
 - Webb, Olver, "Spectra of Jacobi operators via connection coefficient matrices," CIMP, 2021.
 - C., "Computing spectral measures and spectral types," CIMP, 2021.
 - C., Horning, Townsend, "Computing spectral measures of self-adjoint operators," SIAM Rev, 2021.
- Resonances ($\Delta_2 \setminus \Sigma_1 \cup \Pi_1$)
 - Ben-Artzi, Marletta, Rösler, "Computing scattering resonances," JEMS, 2022.
 - Ben-Artzi, Marletta, Rösler, "Computing the sound of the sea in a seashell," FOCM, 2022.
- Zoo of problems (e.g., fractal dims, capacity, radii etc.) up to Σ_6 !!!!
 - C., "On the computation of geometric features of spectra of linear operators on Hilbert spaces," FOCM, under revisions.

General technique for lower bounds by embedding combinatorial problems.



E.g., ground state of quasicrystal



E.g., continuous spectra of graphene

Remarks

- Any model of computation.
 - (lower bds universal, upper bds can be realized via interval arith.)
- Practical: "Spot" info or subclass needed to lower difficulty.
- Tells us what we can or cannot do.
- Existing hierarchies become special cases.
- Extends beyond spectral problems: Foundations of AI, PDEs (e.g., time-dep. Schrödinger eq. on $L^2(\mathbb{R}^d)$ with error control), optimisation (e.g., guarantees), computer-assisted proofs, ...

Example 2: Smale's 18th problem "What are the limits of AI?"

"Very often, the creation of a technological artifact precedes the science that goes with it. The steam engine was invented before thermodynamics. Thermodynamics was invented to explain the steam engine, essentially the limitations of it. What we are after is the equivalent of thermodynamics for intelligence." Yann LeCun

Lower bounds: Use SCI techniques to prove impossibility results. (Different techniques needed for training problems.)

*S. Smale's list of problems for the 21st century (requested by V. Arnold), inspired by Hilbert's list

Problem: hallucinations and instability



"AI generated hallucination", from Facebook and NYU's FastMRI challenge 2020

Instabilities in medical diagnosisOriginal MolePerturbed Mole



Model confidence

Model confidence

From Finlayson et al., "Adversarial attacks on medical machine learning," Science, 2019.

When can we make AI robust and trustworthy?

Example of the limits of deep learning

Paradox: "Nice" linear inverse problems where a *stable* and *accurate* neural network for image reconstruction <u>exists</u>, but it <u>can never be trained</u>!

E.g., suppose we want to solve (holds for much more general problems)

$$\min_{x \in \mathbb{C}^N} \|x\|_{l^1} + \lambda \|Ax - y\|_{l_2}^2$$

 $A \in \mathbb{C}^{m \times N}$ (modality, m < N), $S = \{y_j\}_{j=1}^R$ (samples)

Arises when given $y \approx Ax + e$.

Allow arbitrary precision of training data.

Enforce condition numbers bounded by 1.

Example of the limits of deep learning

Paradox: "Nice" linear inverse problems where a *stable* and *accurate* neural network for image reconstruction <u>exists</u>, but it <u>can never be trained</u>!

Theorem: Pick positive integers $n \geq 3$ and M. Class of problems such that:

- (Not trainable) No algorithm (even random) can train a neural network with n digits of accuracy over the dataset with probability greater than 1/2.
- (Not practical) n 1 digits of accuracy possible over the dataset, but any training algorithm requires arbitrarily large training data.
- (Trainable and practical) n 2 digits of accuracy possible over the dataset via training algorithm using *M* training data.

Holds for any architecture, any precision of training data.

\Rightarrow Classification theory telling us what can and cannot be done

C., Antun, Hansen, "The difficulty of computing stable and accurate neural networks: On the barriers of deep learning and Smale's 18th problem," PNAS, 2022.

- Antun, C., Hansen, "Proving Existence Is Not Enough: : Mathematical Paradoxes Unravel the Limits of Neural Networks in Artificial Intelligence," SIAM News, May 2022.
- Choi, "Some AI Systems May Be Impossible to Compute," IEEE Spectrum, March 2022.

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Phase transitions



SCI machinery for embedding into well-conditioned problems

Idea of mechanism

General lemma (works in other scenarios)



The world of neural networks



Given a problem and conditions, where does it sit in this diagram?

The world of neural networks



Given a problem and conditions, where does it sit in this diagram?

Example counterpart theorem

Certain conditions: <u>stable</u> neural networks <u>trained</u> with <u>exponential accuracy</u>. E.g., *approximate Łojasiewicz-type inequality*:

> (1) $\min_{x \in \mathbb{C}^N} f(x)$ s.t. $||Ax - y|| \le \varepsilon$ dist(x, solution) $\le \alpha([f(x) - f^*] + [||Ax - y|| - \varepsilon] + \delta)$

Fast Iterative **RE**started **NET**works (FIRENETs) (unrolled primal-dual with novel restart scheme)

Theorem: Training algorithm that, under above assumption, produces *stable* neural networks φ_n of width O(N), depth O(n), guaranteed worst bound

 $dist(\varphi_n(y), solution) \leq e^{-n} + \delta$

C., Antun, Hansen, "The difficulty of computing stable and accurate neural networks: On the barriers of deep learning and Smale's 18th problem," PNAS, 2022.

[•] C., "WARPd: A linearly convergent first-order method for inverse problems with approximate sharpness conditions," SIIMS, to appear.

Example of severe instability



• Zhu et al., "Image reconstruction by domain-transform manifold learning," Nature, 2018.

MRI: discrete 2D

• Antun et al., "On instabilities of deep learning in image reconstruction and the potential costs of AI," PNAS, 2020.

^{29/31} FIRENET: provably stable (even to adversarial examples) and accurate



• C., Antun, Hansen, "The difficulty of computing stable and accurate neural networks: On the barriers of deep learning and Smale's 18th problem," PNAS, 2022.

Key pillars: stability and accuracy



• C., Antun, Hansen, "The difficulty of computing stable and accurate neural networks: On the barriers of deep learning and Smale's 18th problem," PNAS, 2022.

U-Net with no noise: accurate but unstable



• C., Antun, Hansen, "The difficulty of computing stable and accurate neural networks: On the barriers of deep learning and Smale's 18th problem," PNAS, 2022.

U-Net with noise: stable but inaccurate



• C., Antun, Hansen, "The difficulty of computing stable and accurate neural networks: On the barriers of deep learning and Smale's 18th problem," PNAS, 2022.

FIRENET: balances stability and accuracy?



• C., Antun, Hansen, "The difficulty of computing stable and accurate neural networks: On the barriers of deep learning and Smale's 18th problem," PNAS, 2022.

FIRENET: balances stability and accuracy?

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Open problem: use the toolkit to precisely prove theorems about *optimal* trade-offs.



• C., Antun, Hansen, "The difficulty of computing stable and accurate neural networks: On the barriers of deep learning and Smale's 18th problem," PNAS, 2022.

Summary: When can algorithms be trusted?

Toolkit + programme:

- Classifying the difficulty of computational problems
- Proving the optimality of algorithms.

Two **examples**:

- Rigorous and verifiable data-driven dynamical systems: use the residual!
- Computational boundaries in AI: we need foundations!

Building blocks for further problems in data analysis and beyond ...

Additional slides

Example of "too much" (spectral pollution)



• C., Roman, Hansen, "How to compute spectra with error control," Physical Review Letters, 2019.

Convergence of quadrature

E.g.,
$$\langle \mathcal{K}\psi_k, \psi_j \rangle = \lim_{M \to \infty} \sum_{m=1}^M w_m \overline{\psi_j(x^{(m)})} \underbrace{\psi_k(y^{(m)})}_{[\mathcal{K}\psi_k](x^{(m)})}$$

Three examples:

- **High-order quadrature:** $\{x^{(m)}, w_m\}_{m=1}^{M} M$ -point quadrature rule. Rapid convergence. Requires free choice of $\{x^{(m)}\}_{m=1}^{M}$ and small d.
- Random sampling: $\{x^{(m)}\}_{m=1}^{M}$ selected at random. Most common Large *d*. Slow Monte Carlo $O(M^{-1/2})$ rate of convergence.
- Ergodic sampling: $x^{(m+1)} = F(x^{(m)})$. Single trajectory, large d. Requires ergodicity, convergence can be slow.



C., Townsend, "Rigorous data-driven computation of spectral properties of Koopman operators for dynamical systems,"
 Communications on Pure and Applied Mathematics, under review.

Example: wall-jet boundary layer

- b
- 12,000 snapshots over 1s
- Reynolds number $\approx 6.4 \times 10^4$
- Ambient dimension ≈ 100,000 (velocity at measurement points)

*Raw measurements provided by Máté Szőke (Virginia Tech)









Koopman mode decomposition ($\mathbb{K}V = V\Lambda$)

Standard Koopman mode decomposition (order modes by $|\Lambda|$):

$$g(x) \approx \underbrace{\left[\psi_{1}(x) \cdots \psi_{N_{K}}(x)\right]V}_{\text{approx Koopman e-functions}} \underbrace{\left(V\sqrt{W}\Psi_{X}\right)^{\dagger}\sqrt{W}\left[g(x^{(1)}) \cdots g(x^{(M)})\right]^{T}}_{\text{Koopman modes}}$$

$$\stackrel{?}{\Rightarrow} g(x_{n}) \approx \underbrace{\left[\psi_{1}(x) \cdots \psi_{N_{K}}(x)\right]V}_{\text{approx Koopman e-functions}} \Lambda^{n} \underbrace{\left(V\sqrt{W}\Psi_{X}\right)^{\dagger}\sqrt{W}\left[g(x^{(1)}) \cdots g(x^{(M)})\right]^{T}}_{\text{Koopman modes}}$$
Residual Koopman mode decomposition (order modes by res (λ, \mathbf{v})):
$$g(x) \approx \underbrace{\left[\psi_{1}(x) \cdots \psi_{N_{K}}(x)\right]V_{(\mathcal{E})}}_{\text{approx Koopman e-functions}} \underbrace{\left(V_{(\mathcal{E})}\sqrt{W}\Psi_{X}\right)^{\dagger}\sqrt{W}\left[g(x^{(1)}) \cdots g(x^{(M)})\right]^{T}}_{\text{Koopman modes}}$$

$$g(x_{n}) \approx \underbrace{\left[\psi_{1}(x_{0}) \cdots \psi_{N_{K}}(x_{0})\right]V_{(\mathcal{E})}}_{\text{approx Koopman e-functions}} \Lambda^{n}_{(\mathcal{E})} \underbrace{\left(V_{(\mathcal{E})}\sqrt{W}\Psi_{X}\right)^{\dagger}\sqrt{W}\left[g(x^{(1)}) \cdots g(x^{(M)})\right]^{T}}_{\text{Koopman modes}}$$

Example: laser-induced plasma



Example: ionization probabilities in < 1 second

$$[Au](r) = -\frac{d^2u}{dr^2}(r) + \left(\frac{l(l+1)}{r^2} + \frac{e^{-r} - 1}{r}\right)u(r), \qquad r > 0$$

Hellman potential

% Norm squared

- % Measure wrt g(r)
- % Potential, 1=1
- % Quadrature rule
- % epsilon=0.1, m=4
- % Ionization prob

Example: ionization probabilities in < 1 second



Extensions to PDEs

• Operator A on \mathbb{R}^d of form

$$[Au] = \sum_{k \in \mathbb{Z}^d_{\geq 0}, |k| \leq N} c_k(x) [\partial^k u](x)$$

- Assume coefficients are
 - Polynomially bounded
 - Locally bounded total variation
- Build matrix representation using quasi-Monte Carlo integration.
- \rightarrow Sample coefficients to compute Spec(A) with error control!
- **NB:** works for spectral methods, FEM, etc.

Eigenvalues of Dirac operator



Example of impossibility theorem for subsampled discrete cosine transform

$\min_k \operatorname{dist}(\Psi_n(y_k), \Xi(A, y_k))$	$\min_k \operatorname{dist}(\Phi_n(y_k), \Xi(A, y_k))$	prec. of training data	$10^{-n^{-1}}$
0.2999690	0.2597827	m = 10	10^{-1}
0.3000000	0.2598050	m = 20	10^{-1}
0.3000000	0.2598052	m = 30	10^{-1}
0.0030000	0.0025980	m = 10	10^{-3}
0.0030000	0.0025980	m = 20	10^{-3}
0.0030000	0.0025980	m = 30	10^{-3}
0.0000030	0.000015	m = 10	10^{-6}
0.0000030	0.000015	m = 20	10^{-6}
0.0000030	0.000015	$m \equiv 30$	10^{-6}

Impossibility of computing approximations of the neural network to arbitrary accuracy. We demonstrate the impossibility statement on fast iterative restarted networks Φ_n and learned iterative shrinkage thresholding algorithm networks Ψ_n . The table reveals the shortest l_2 distance between the networks' output and the problem's true solution for different values of m (precision of training data is 2^{-m}). Neither of the trained neural networks can compute the existing correct neural network to accuracy 10^{-n} , but both compute approximations that are accurate to 10^{-n+1} .

Stabilising unstable neural networks

