# Infinite dimensional spectral computations \& linear algebra: 

Extending the QR algorithm to infinite dimensions

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## Outline

- Background
- Introducing IQR
- Non Normal Operators
- How To Compute
- Numerical Examples
- Conclusion

Background

## Background

- Hilbert space $I^{2}(\mathbb{N})$ with $\|x\|_{2}=\sqrt{\sum_{j=1}^{\infty}}\left|x_{j}\right|^{2},\langle x, y\rangle=\sum_{j=1}^{\infty} x_{j} \bar{y}_{j}$
- Bounded linear operator $T: I^{2}(\mathbb{N}) \rightarrow I^{2}(\mathbb{N})$ realised as matrix

$$
\left(\begin{array}{cccc}
t_{11} & t_{12} & t_{13} & \ldots \\
t_{21} & t_{22} & t_{23} & \ldots \\
t_{31} & t_{32} & t_{33} & \ldots \\
\vdots & \vdots & \vdots & \ddots
\end{array}\right)
$$

Denote these by $\mathcal{B}\left(I^{2}(\mathbb{N})\right)$.

- Want to compute spectrum (generalistion of eigenvalues)

$$
\sigma(T):=\{z \in \mathbb{C}: T-z l \text { not invertible }\} .
$$

from the matrix elements. What about eigenvectors etc.?

## Well Studied

- Quantum mechanics, quasicrystals


Figure: Left: Dan Shechtman, Nobel Prize in Chemistry 2011. Right: Electron diffraction pattern of quasicrystal.

- Intensely investigated since the 1950s, still very active today.


Figure: Left: Artur Avila, Fields Medal 2014. Right: Hofstadter butterfly.

## Hierarchy of complexity

Definition (Tower of Algorithms)
A tower of algorithms of height $k$ is a family of sequences of functions

$$
\Gamma_{n_{k}, \ldots, n_{1}}: \Omega \rightarrow \mathcal{M}
$$

where $n_{k}, \ldots, n_{1} \in \mathbb{N}$ and $\Gamma_{n_{k}, \ldots, n_{1}}$ are "algorithms". Moreover,

$$
\sigma(T)=\lim _{n_{k} \rightarrow \infty} \ldots \lim _{n_{1} \rightarrow \infty} \Gamma_{n_{k}, \ldots, n_{1}}(T)
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## Definition (Solvability Complexity Index (SCI))

Solvability Complexity $\operatorname{Index}, \operatorname{SCI}(\sigma, \Omega)$ is the smallest integer $k$ for which there exists a tower of algorithms of height $k$. If no such tower exists then $\operatorname{SCI}(\sigma, \Omega)=\infty$.

## Example Results

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(1) Splitting the discrete spectrum from essential spectrum (as sets) in generally harder.

Methods based on approximating pseudospectrum:

$$
\sigma_{\epsilon}(T)=\left\{z:\left\|(T-z l)^{-1}\right\| \geq \epsilon^{-1}\right\}
$$

where we interpret $\left\|S^{-1}\right\|$ as $+\infty$ if $S$ does not have a bounded inverse.

## Motivation

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## Motivation

(1) Above method can't detect isolated eigenvalues and their multiplicity.
(2) Due to taking square root, above method can only gain precision $\sqrt{\epsilon_{\text {mach }}}$.
(3) Can we generalise staple finite matrix algorithms to infinite dimensions?
(1) Can we gain error control and classification results in the hierarchy?

## Classical QR

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T=Q_{1} R_{1}
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& \vdots \\
T_{n} & =Q_{n}^{*} \ldots Q_{1}^{*} T Q_{1} \ldots Q_{n}
\end{aligned}
$$

## Classical Result

## Theorem

Let $T \in \mathbb{C}^{N \times N}$ be a normal matrix with eigenvalues satisfying $\left|\lambda_{1}\right|>\ldots>\left|\lambda_{N}\right|$. Let $\left\{Q_{m}\right\}$ be a $Q$-sequence of unitary operators. Then (up to re-ordering of the basis)

$$
Q_{m}^{*} T Q_{m} \longrightarrow \bigoplus_{j=1}^{N} \lambda_{j} e_{j} \otimes e_{j}, \quad m \rightarrow \infty
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Numerical Example ...

Introducing IQR

## The Main Idea

\[

\]

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Iterate QR? Truncate

$$
\left(\begin{array}{cccc}
t_{11} & t_{12} & t_{13} & \cdots \\
t_{21} & t_{22} & t_{23} & \cdots \\
\vdots & \vdots & \vdots & \ddots
\end{array}\right) \Rightarrow\left(\begin{array}{ccc}
\tilde{d}_{11} & 0 & \ldots \\
0 & \tilde{d}_{22} & \cdots \\
\vdots & \vdots & \ddots
\end{array}\right) \Rightarrow\left(\begin{array}{cccc}
\tilde{d}_{11} & 0 & \ldots & 0 \\
0 & \tilde{d}_{22} & \cdots & 0 \\
\vdots & \vdots & \ddots & \vdots \\
0 & 0 & \cdots & \tilde{d}_{n n}
\end{array}\right)
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## Questions

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(1) Can we even compute this beast on a finite machine?

## The QR Decomposition

Definition
A Householder reflection is an operator $S \in \mathcal{B}(\mathcal{H})$ of the form

$$
S=I-\frac{2}{\|\psi\|^{2}} \psi \otimes \bar{\psi}, \quad \psi \in \mathcal{H}
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where $\bar{\psi}$ denotes the associated functional in $\mathcal{H}^{*}$ given by $x \rightarrow\langle x, \psi\rangle$.

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where $\bar{\psi}$ denotes the associated functional in $\mathcal{H}^{*}$ given by $x \rightarrow\langle x, \psi\rangle$.In the case where $\mathcal{H}=\mathcal{H}_{1} \oplus \mathcal{H}_{2}$ and $\boldsymbol{I}_{i}$ is the identity on $\mathcal{H}_{i}$ then

$$
U=I_{1} \oplus\left(I_{2}-\frac{2}{\|\psi\|^{2}} \psi \otimes \bar{\psi}\right) \quad \psi \in \mathcal{H}_{2}
$$

is called a Householder transformation.

## The QR Decomposition

Theorem (Hansen 2008)
Let $T$ be a bounded operator on a separable Hilbert space $\mathcal{H}$ and let $\left\{e_{j}\right\}_{j \in \mathbb{N}}$ be an orthonormal basis for $\mathcal{H} \cong I^{2}(\mathbb{N})$. Then there exist an isometry $Q$ such that $T=Q R$ where $R$ is upper triangular with respect to $\left\{e_{j}\right\}$. Moreover,

$$
Q=\underset{n \rightarrow \infty}{\text { SOT-lim }} V_{n}
$$

where $V_{n}=U_{1} \cdots U_{n}$ are unitary and each $U_{j}$ is a Householder transformation.

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Good for numerics - more stable than Gram-Schmidt...

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Assume $T$ invertible...

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An easy induction gives us that

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$$

$\hat{R}_{m}$ upper triangular since $R_{j}, j \leq m$ are. By invertibility of $T$, $\left\langle R e_{i}, e_{i}\right\rangle \neq 0$. Hence

$$
\operatorname{span}\left\{T^{m} e_{j}\right\}_{j=1}^{J}=\operatorname{span}\left\{\hat{Q}_{m} e_{j}\right\}_{j=1}^{J}, \quad J \in \mathbb{N}
$$

## Questions

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## A Result for Normal Operators

Assume the following:
(A1) $T \in \mathcal{B}(\mathcal{H})$ is an invertible normal operator and $\left\{e_{j}\right\}_{j \in \mathbb{N}}$ an orthonormal basis for $\mathcal{H}$. $\left\{Q_{k}\right\}$ and $\left\{R_{k}\right\}$ are $Q$ - and $R$-sequences of $T$ with respect to the basis $\left\{e_{j}\right\}_{j \in \mathbb{N}}$.
(A2) $\sigma(T)=\omega \cup \Psi$ such that $\omega \cap \Psi=\emptyset$ and $\omega=\left\{\lambda_{i}\right\}_{i=1}^{N}$, where the $\lambda_{i} \mathrm{~s}$ are isolated eigenvalues with (possibly infinite) multiplicity $m_{i}$. Let $M=m_{1}+\ldots+m_{N}=\operatorname{dim}\left(\operatorname{ran} \chi_{\omega}(T)\right)$ and suppose that $\left|\lambda_{1}\right|>\ldots>\left|\lambda_{N}\right|$. Suppose further that $\sup \{|\theta|: \theta \in \Psi\}<\left|\lambda_{N}\right|$.

Define

$$
\rho=\sup \{|z|: z \in \Psi\}, \quad r=\max \left\{\left|\lambda_{2} / \lambda_{1}\right|, \ldots,\left|\lambda_{N} / \lambda_{N-1}\right|, \rho /\left|\lambda_{N}\right|\right\}
$$

then $r<1$.

## What This Really Means!



## A Result for Normal Operators

## Theorem

There exists $\left\{\hat{e}_{j}\right\}_{j=1}^{M} \subset\left\{e_{j}\right\}_{j \in \mathbb{N}}$, where $M=m_{1}+\ldots+m_{N}$, so that $\operatorname{span}\left\{Q_{k} \hat{e}_{j}\right\} \rightarrow \operatorname{span}\left\{\hat{q}_{j}\right\}$ where $\left\{\hat{q}_{j}\right\}_{j=1}^{M} \subset \operatorname{ran} \chi_{\omega}(T)$ is a collection of orthonormal eigenvectors of $T$ and if $e_{j} \notin\left\{\hat{e}_{j}\right\}_{j=1}^{M}$, then $\chi_{\omega}(T) Q_{k} e_{j} \rightarrow 0$. Also:
(i) Every subsequence of $\left\{Q_{n}^{*} T Q_{n}\right\}_{n \in \mathbb{N}}$ has a convergent subsequence $\left\{Q_{n_{k}}^{*} T Q_{n_{k}}\right\}_{k \in \mathbb{N}}$ such that

$$
Q_{n_{k}}^{*} T Q_{n_{k}} \xrightarrow{\text { WOT }}\left(\bigoplus_{j=1}^{M}\left\langle T \hat{q}_{j}, \hat{q}_{j}\right\rangle \hat{e}_{j} \otimes \hat{e}_{j}\right) \oplus \sum_{j \in \Theta} \xi_{j} \otimes e_{j},
$$

as $k \rightarrow \infty$, where

$$
\Theta=\left\{j: e_{j} \notin\left\{\hat{e}_{j}\right\}_{j=1}^{M}\right\}, \quad \xi_{j} \in \overline{\operatorname{span}\left\{e_{i}\right\}_{i \in \Theta}}
$$

and only $\sum_{j \in \Theta} \xi_{j} \otimes e_{j}$ depends on the choice of subsequence.
(ii)

$$
\widehat{P}_{M} Q_{n}^{*} T Q_{n} \widehat{P}_{M} \xrightarrow{\operatorname{SOT}}\left(\bigoplus_{j=1}^{M}\left\langle T \hat{q}_{j}, \hat{q}_{j}\right\rangle \hat{e}_{j} \otimes \hat{e}_{j}\right), \quad \text { as } n \rightarrow \infty,
$$

where $\widehat{P}_{M}$ denotes the orthogonal projection onto $\overline{\operatorname{span}\left\{\hat{e}_{j}\right\}_{j=1}^{M}}$. For any fixed $x \in \operatorname{span}\left\{\hat{e}_{j}\right\}_{j=1}^{M}$ we have the following rate of convergence

$$
\left\|\widehat{P}_{M} Q_{n}^{*} T Q_{n} \widehat{P}_{M^{x}}-\left(\bigoplus_{j=1}^{M}\left\langle T \hat{q}_{j}, \hat{q}_{j}\right\rangle \hat{e}_{j} \otimes \hat{e}_{j}\right) \times\right\| \leq \mathcal{O}\left(r^{n}\right)
$$

## Idea of Proof

If we keep applying $T$ to $\operatorname{span}\left\{e_{j}\right\}_{j=1}^{K}$, expect the span of vectors will approach the extreme parts of spectrum (projection sense).

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where $\left\{\xi_{j}\right\}_{j=1}^{M}$ is an orthonormal set of eigenvectors of $T$ to make this precise.

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Finally use

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$$

Can upgrade for block convergence (eigenvalues not of distinct magnitude), SOT convergence etc.

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# Non Normal Operators 

## Set Up

Assume the following ( $M$ now finite):
(A1) $T \in \mathcal{B}(\mathcal{H})$ is an invertible operator and there is an orthogonal projection $P$ of finite rank $M$ with image invariant under $T$.
(A2) There exist $\alpha>\beta>0$ such that

$$
\begin{aligned}
& \|T x\| \geq \alpha\|x\| \quad \forall x \in \operatorname{ran}(P) \\
& \|(I-P) T(I-P)\| \leq \beta
\end{aligned}
$$

(A3) $\left\{\tilde{P} e_{j}\right\}_{j=1}^{M}$ are linearly independent ( $\tilde{P}$ canoncal defined later).
Define for orthogonal projections $E, F$ onto $S_{1}, S_{2} \subset \mathcal{H}$ the distance between subspaces

$$
\hat{\delta}\left(S_{1}, S_{2}\right)=\|E-F\| \in[0,1]
$$

and the subspace angle

$$
\phi\left(S_{1}, S_{2}\right)=\sin ^{-1}\left(\hat{\delta}\left(S_{1}, S_{2}\right)\right) .
$$

## Result

## Theorem

There exists a canonical $M$ dimensional $T^{*}$-invariant subspace $S$ with orthogonal projection $\tilde{P}$ with
(i) The subspace angle $\phi\left(\operatorname{span}\left\{e_{j}\right\}_{j=1}^{M}, S\right)<\pi / 2$ and we have
$\hat{\delta}\left(\operatorname{span}\left\{Q_{n} e_{j}\right\}_{j=1}^{M}, \operatorname{ran}(P)\right) \leq \frac{\sin \left(\phi\left(\operatorname{span}\left\{e_{j}\right\}_{j=1}^{M}, \operatorname{ran}(P)\right)\right)}{\cos \left(\phi\left(\operatorname{span}\left\{e_{j}\right\}_{j=1}^{M}, S\right)\right)}\left(1+\frac{\|P T(I-P)\|}{\alpha-\beta}\right) \frac{\beta^{n}}{\alpha^{n}}$,
(ii) Every subsequence of $\left\{Q_{n}^{*} T Q_{n}\right\}_{n \in \mathbb{N}}$ has a convergent subsequence $\left\{Q_{n_{k}}^{*} T Q_{n_{k}}\right\}_{k \in \mathbb{N}}$ such that

$$
\begin{gathered}
Q_{n_{k}}^{*} T Q_{n_{k}} \xrightarrow{\text { wOT }} \sum_{j=1}^{M} \xi_{j} \otimes e_{j} \bigoplus \sum_{i=M+1}^{\infty} \zeta_{i} \otimes e_{i}, \\
\xi_{j} \in \overline{\operatorname{span}\left\{e_{j}\right\}_{j=1}^{M}}, \quad \zeta_{i} \in \mathcal{H} .
\end{gathered}
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How To Compute

## Some Definition

## Definition

Let $T$ be an infinite matrix acting as a bounded operator on $I^{2}(\mathbb{N})$ with basis $\left\{e_{j}\right\}_{j \in \mathbb{N}}$. For $f: \mathbb{N} \rightarrow \mathbb{N}$ non-decreasing with $f(n) \geq n$ we say that $T$ has quasi-banded subdiagonals with respect to $f$ if $\left\langle T e_{j}, e_{i}\right\rangle=0$ when $i>f(j)$.

## A Theorem

## Theorem

Let $T \in \mathcal{B}\left(I^{2}(\mathbb{N})\right)$ have quasi-banded subdiagonals with respect to $f$ and let $T_{n}$ be the $n$-th element in the QR iteration, i.e.
$T_{n}=Q_{n}^{*} \cdots Q_{1}^{*} T Q_{1} \cdots Q_{n}$, where

$$
Q_{j}=\underset{I \rightarrow \infty}{\operatorname{SOT}-\lim } U_{1}^{j} \cdots U_{I}^{j}
$$

and $U_{l}^{j}$ is a Householder transformation. Let $P_{m}$ be the usual projection onto $\operatorname{span}\left\{e_{j}\right\}_{j=1}^{m}$ and denote the a-fold iteration of $f$ by $\underbrace{f \circ f \circ \ldots \circ f}_{\text {a times }}=f_{a}$. Then
a times
$P_{m} T_{n} P_{m}=P_{m} U_{m}^{n} \cdots U_{1}^{n} U_{f_{1}(m)}^{n-1} \cdots U_{1}^{n-1} \cdots U_{f_{(n-2)}(m)}^{2} \cdots U_{1}^{2} U_{f_{(n-1)}(m)}^{1} \cdots U_{1}^{1}$

- $P_{f_{n}(m)} T P_{f_{n}(m)}$
$\cdot U_{1}^{1} \cdots U_{f_{(n-1)}(m)}^{1} U_{1}^{2} \cdots U_{f_{(n-2)}(m)}^{2} \cdots U_{1}^{n-1} \cdots U_{f_{1}(m)}^{n-1} U_{1}^{n} \cdots U_{m}^{n} P_{m}$


## Why?

In the subcase of invertibility, a consequence of the fact that if $T$ has quasi-banded subdiagonals with respect to $f$ then

$$
P_{m} T^{n} P_{m}=P_{m}\left(P_{f_{n}(m)} T P_{f_{n}(m)}\right)^{n} P_{m}
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## Why?

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We can apply Gram-Schmidt (or a more stable modified version) to the columns of $P_{f_{n}(m)} T P_{f_{n}(m)}$ and truncate the resulting matrix!

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$$

We can apply Gram-Schmidt (or a more stable modified version) to the columns of $P_{f_{n}(m)} T P_{f_{n}(m)}$ and truncate the resulting matrix!
Can also extend to compute the IQR iterates with error control if we can evaluate an increasing family of increasing functions $g^{j}: \mathbb{N} \rightarrow \mathbb{N}$ such that defining the matrix $T_{(j)}$ with columns $\left\{P_{g^{j}(n)} T e_{n}\right\}$ we have that $T_{(j)}$ is invertible and

$$
\left\|\left(P_{g^{j}(n)}-l\right) T e_{n}\right\| \leq \frac{1}{j} .
$$

## Questions

© Does the QR algorithm exist in infinite dimensions?
(2) When do we gain convergence to a diagonal operator and in what sense?
(3) Can we prove anything for non normal operators in infinite dimensions?
(1) Can we even compute this beast on a finite machine?

# Numerical Examples 

## Example 1

Take unilateral shift $U: e_{n} \rightarrow e_{n+1}$ acting on $I^{2}(\mathbb{Z})$ (with the natural choice of indexing $\mathbb{Z}$ ) and perturb by the compact diagonal operator

$$
D\left(e_{n}\right)=\frac{5 \sin (n)^{2}}{\sqrt{|n|+1}} e_{n} .
$$

Hence the spectrum of the full operator $T=U+D$ consists of the unit circle and a collection of eigenvalues in the discrete spectrum.

## Example 1



Figure: Left: Error in approximating $\lambda_{1}$ (top) and $\lambda_{2}$ through taking finite sections $P_{100} Q_{n}^{*} T Q_{n} \mid P_{100} \mathcal{H}$. We have shown the expected rates of convergence as references. Right: The spectrum of $\left.P_{100} Q_{500}^{*} T Q_{500}\right|_{P_{100} \mathcal{H}}$ demonstrating convergence to the extremal parts.

## Example 2

Almost Mathieu related to a wealth of mathematical/physical problems:

$$
\left(H_{1} x\right)_{n}=x_{n-1}+x_{n+1}+2 \cos (2 \pi n \alpha+\nu) x_{n}
$$

on $I^{2}(\mathbb{Z})$. Hamiltonian represents crystal electron in a uniform magnetic field and the spectrum the allowed energies of the system. No discrete spectrum!

## Example 2



Figure: The spectrum of $H_{1}$ calculated analytically for rational $\alpha$ and the output of finite section for $m=100$. Note that the finite section method causes strong spectral pollution.

## Example 2




Figure: IQR for $m=100, n=100$ and the same for $m=100, n=5000$. IQR algorithm is more robust, preserving the structure of the spectrum whilst converging to the boundary of the essential spectrum.

## Example 3

Toeplitz operator

$$
N=\frac{1}{2}\left(U_{3}+U_{-1}\right)
$$

where $U_{m}$ acts on $I^{2}(\mathbb{Z})$ by the shift $e_{j} \rightarrow e_{j+m}$ on the standard basis. Not invertible and has no eigenvalues. With no shift, $Q_{n}^{*} N Q_{n}$ appeared to converge strongly to the operator

$$
\tilde{N}=\left(\begin{array}{lll}
A & & \\
& A & \\
& & \ddots
\end{array}\right), \quad A=\left(\begin{array}{llll}
0 & 1 & 0 & 0 \\
0 & 0 & 0 & 1 \\
1 & 0 & 0 & 0 \\
0 & 0 & 1 & 0
\end{array}\right)
$$

Spectrum equal to $\{ \pm 1, \pm 1 i\}$ which are the extremal points of $\sigma(N)$. Shift to $Q_{n}^{*}(N+I) Q_{n}$, then we appeared to converge to the diagonal operator $D=2 I$. BUT curious case of reduced rate of convergence...

## Example 3



Figure: Left: Spectrum of the normal operator $N$ and finite section approximates. Right: Convergence of first five diagonal entries of $Q_{n}^{*}(N+I) Q_{n}$ to 2. Convergence of the rate $\mathcal{O}(1 / n)$ for this part of the essential spectrum as opposed to the linear convergence rate $\mathcal{O}\left(r^{n}\right)$ seen in the first example.

## Example 4

$$
A=\left(\begin{array}{ccccccc}
0 & 0 & 0 & 0 & 0 & 0 & \ldots \\
1 & 0 & a_{23} & a_{24} & 0 & 0 & \ldots \\
0 & 1 & 0 & 0 & 0 & 0 & \ldots \\
0 & 0 & 1 & 0 & a_{45} & a_{46} & \ldots \\
0 & 0 & 0 & 1 & 0 & 0 & \ldots \\
0 & 0 & 0 & 0 & 1 & 0 & \ldots \\
\vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots
\end{array}\right),
$$

where $a_{2 j, 2 j+1}=i$, and $a_{2 j, 2 j+2}=-i$ if $j$ is prime and $a_{2 j, 2 j+2}=0$ otherwise.

## Example 4

$$
T=\left(\begin{array}{cccccccc}
2.5+0.5 i & 0 & 0 & 0 & 0 & 0 & 0 & \ldots \\
1 & 3-0.5 i & 0 & 0 & 0 & 0 & 0 & \ldots \\
0 & 1 & 1.7 & 0.05 & 0 & 0 & 0 & \ldots \\
0 & 0 & 0.05 & t_{4} & 0 & 0 & 0 & \ldots \\
0 & 0 & 0 & 0 & t_{5} & 0 & 0 & \ldots \\
0 & 0 & 0 & 0 & 1 & t_{6} & 0 & \ldots \\
0 & 0 & 0 & 0 & 0 & 1 & t_{7} & \ldots \\
\vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots
\end{array}\right)
$$

where $t_{j}=1+0.5(\sin (j)+i \cos (j))$ for $j \geq 4$.

## Example 4



Figure: Left: $\sigma\left(\left.P_{n} A\right|_{P_{n} \mathcal{H}}\right)$ for $n=1000$ with the false eigenvalue (recall that $\left.\sigma(A) \subseteq \sigma_{\epsilon}(A)\right)$. Right: $\sigma\left(\left.P_{n} T\right|_{P_{n} \mathcal{H}}\right)$ for $n=500$ along with contours of the resolvent norm and $\sigma_{\epsilon}(T)$ for $\epsilon=2 \sqrt{\epsilon_{\text {mach }}}$.

## Example 4




Figure: The figures show $\sigma\left(P_{m} Q_{n}^{*} A Q_{n} \mid P_{m} \mathcal{H}\right)$ (left), $\sigma\left(P_{m} Q_{n}^{*} T Q_{n} \mid P_{m} \mathcal{H}\right)$ (right) for $n=1000, m=1000$ and $n=1500, m=500$ respectively.

## Example 4




Figure: Left: Absolute values of entries of $Q_{1000}^{*} A Q_{1000}$. Right: Absolute values of entries of $Q_{1500}^{*} T Q_{1500}$. We appear to gain convergence to a diagonally dominated upper triangular matrix.

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(1) Paper available soon!

## Future Work

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(3) Many more open computational spectral problems - spectral measures, IQR for unbounded operators etc.

## References

Only two main references:
(1) First appearance of the algorithm for real symmetric infinite matrices: P Deift, LC Li, and C Tomei. Toda flows with infinitely many variables. Journal of functional analysis, 64(3):358402, 1985.
(2) Existence of QR decomposition for non invertible operators and eigenvector convergence theorem for normal operators: AC Hansen. On the approximation of spectra of linear operators on Hilbert spaces. J. Funct. Anal., 254(8):20922126, 2008.
Other SCI results: In progress!

Thanks for listening!
Any questions?

"Wouldn't it be more efficient to just find who's complicating equations and ask them to stop?"

