## Exercise sheet 2

- 1. Let  $f : \mathbb{R}^n \to \mathbb{R}$  convex. Show that f is G-Lipschitz (with respect to the  $\ell_2$  norm) iff  $||g||_2 \leq G$  for all  $g \in \partial f(x)$  for all  $x \in \mathbb{R}^n$ .
- 2. Polyak step size for subgradient method: show that the subgradient method with step size  $t_i = (f(x_i) f^*)/||g_i||_2^2$  gives iterates  $f_{\text{best},k}$  that converge to  $f^*$  at the rate  $1/\sqrt{k}$  (hint: start from the inequalities relating  $||x_{k+1} x^*||_2^2$  to  $||x_k x^*||_2^2$ ).
- 3. To minimize a nonsmooth function f on a convex set C, the projected subgradient method proceeds as follows:  $x_{k+1} = P_C(x_k t_k g_k)$  where  $P_C$  is the Euclidean projection on C, and  $g_k \in \partial f(x_k)$ . Analyze the convergence of the projected subgradient descent.
- 4. Compute conjugates of following functions
  - (a)  $f(x) = \sum_{i=1}^{n} x_i \log x_i$
  - (b)  $f(x) = -\sum_{i=1}^{n} \log x_i$
  - (c)  $f(X) = -\log \det X$  where X is an  $n \times n$  real symmetric positive definite matrix
- 5. Show that for any closed convex function f we have  $\operatorname{prox}_{f}(x) + \operatorname{prox}_{f^*}(x) = x$  for all x (Moreau's identity).
- 6. Let  $f(x) = \max_{i=1}^{m} a_i^T x + b_i$ . Show that  $f(x) = h^*(Ax + b)$  where h(y) is the indicator function of the simplex  $\{y \in \mathbb{R}^m : y_i \ge 0 \ \forall i = 1, \dots, m \text{ and } \sum y_i = 1\}$ . Find an expression for  $f_{\mu}(x) = (h + \mu d)^*(Ax + b)$  where  $d(y) = \sum_{i=1}^{m} y_i \log y_i \log m$ .
- 7. Implement the subgradient method to minimize  $||Ax b||_1$  where A and b are generated at random. Experiment with different choices of step size. Compare with the smoothing method of Nesterov.
- 8. In this exercise we prove a lower complexity bound for nonsmooth convex optimization. Consider an algorithm that starts at  $x_0 = 0$  and such that when applied to a function f, the (i + 1)'th iterate satisfies

$$x_i \in \operatorname{span} \{g_0, \dots, g_i\} \tag{1}$$

where  $g_0 \in \partial f(x_0) = \partial f(0), \dots, g_i \in \partial f(x_i)$ .

(a) Consider the function

$$f(x) = \max_{i=1,\dots,n} x_i + \frac{1}{2} \|x\|_2^2$$

with  $x \in \mathbb{R}^n$ . Compute  $\partial f(x)$  for any x.

- (b) Compute  $f^* = \min_{x \in \mathbb{R}^n} f(x)$  and find a minimizer  $x^*$ .
- (c) Show that f is (1 + R)-Lipschitz on the Euclidean ball  $\{x \in \mathbb{R}^n : ||x||_2 \leq R\}$  [Hint: consider  $||g||_2$  for  $g \in \partial f(x)$ .]
- (d) A first-order oracle for f gives, for any  $x \in \mathbb{R}^n$ , an element  $g \in \partial f(x)$ . Show that one can design a specific first-order oracle for f ensuring that  $x_i$  satisfying (1) is always supported on the first i components only (i.e., the components  $i + 1, \ldots, n$  are zero).

(e) Set n = k + 1. Show that for any algorithm satisfying (1), the following holds:

$$\frac{f_{\text{best},k} - f^*}{G \|x_0 - x^*\|_2} \ge \frac{c}{\sqrt{k+1}}$$

for a constant c > 0, where  $f_{\text{best},k} = \min\{f(x_0), \ldots, f(x_k)\}$  and G is the Lipschitz constant of f on the Euclidean ball of radius  $||x_0 - x^*||_2$ .