Quadrature Problems for Stochastic Differential Equations

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OUTLINE

- I. Computational Problems for SDEs
- II. Deterministic Quadrature on the Lipschitz Class
- III. Randomized Quadrature on the Lipschitz Class
- IV. Quadrature on the Sequence Space

Joint work with

- J. Creutzig (Darmstadt), S. Dereich (Marburg),
- F. Hickernell (IIT Chicago), S. Mayer (Darmstadt),
- T. Müller-Gronbach (Passau), Ben Niu (IIT Chicago),
- S. Toussaint (Darmstadt), L. Yaroslavtseva (Passau).

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SDE

$$dX_t = \mu(X_t) dt + \sigma(X_t) dW_t, \qquad t \in [0, T],$$
$$X_0 = x_0$$

with a Brownian motion W. Solution $X = (X_t)_{t \in [0,T]}$ is a stochastic process with continuous paths,

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Computational problems

- 1. Strong approximation: approximate the solution X.
- 2. Weak approximation: approximate the distribution P_X of X on \mathfrak{X} .
- 3. Quadrature: approximate integrals $E(f(X)) = \int_{\mathfrak{X}} f \, dP_X$.

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Analogously, for the solution X_T at time T.

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Reasonable, but not mandatory, $1. \rightsquigarrow 2. \rightsquigarrow 3$.

$$I(f) = \int_{\mathfrak{X}} f \, dP_X.$$

Deterministic quadrature formulas

$$Q_n(f) = \sum_{i=1}^n a_i \cdot f(x_i)$$

with $a_i \in \mathbb{R}$ and $x_i \in \mathfrak{X}$.

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Maximal error of Q_n on a class F of integrands f

$$e(Q_n, F) = \sup_{f \in F} |I(f) - Q_n(f)|.$$

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n-th minimal error

$$e^{\det}(n,F) = \inf_{Q_n} e(Q_n,F).$$

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 $n\text{-th} \min al \ error$

$$e^{\operatorname{ran}}(n,F) = \inf_{Q_n} e(Q_n,F).$$

For quadrature and weak approximation of SDEs

- 1. distribution P_X on \mathfrak{X} given only implicitly,
- 2. $(\mu, \sigma, x_0) \mapsto \int_{\mathfrak{X}} f \, dP_X$ nonlinear,

3. dim $\mathfrak{X} = \infty$.

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Particular issues

- cost for computation of $Q_n(f)$,
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Assumptions

- scalar SDE (for simplicity),
- $\mu, \sigma \in C^2_{\mathrm{b}}(\mathbb{R})$ (smoothness crucial),
- $\sigma(x_0) \neq 0$ (to exclude deterministic equations).

II. Deterministic Quadrature on the Lipschitz Class

Here $F = \operatorname{Lip}(1)$, i.e., for $\mathfrak{X} = C([0,T])$ and $f \in F$

 $|f(x) - f(y)| \le ||x - y||_{\infty}, \qquad x, y \in \mathfrak{X}.$

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Example: f payoff of an asian or lookback option,

$$f(x) = \max\left(\frac{1}{T}\int_0^T x(t) dt - K, 0\right),$$
$$f(x) = \max\left(\sup_{t \in [0,T]} x(t) - K, 0\right).$$

Equivalence of quadrature on $\operatorname{Lip}(1)$ and quantization of P_X

$$e^{\det}(n,\operatorname{Lip}(1)) = \inf_{x_1,\dots,x_n \in \mathfrak{X}} \operatorname{E}(g(X;x_1,\dots,x_n))$$

for every separable Banach space \mathfrak{X} , where

$$g(x; x_1, \dots, x_n) = \min_{i=1,\dots,n} ||x - x_i||.$$

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Quantization for stochastic processes studied since 2000,

Aurzada, Creutzig, Dereich, Fehringer, Graf, Luschgy, Müller-Gronbach, Matoussi, Pagès, Printems, R, Scheutzow, Wilberts, . . .

In particular, Gaussian processes, Lévy processes, SDEs.

For applications in finance see Pagès, Printems (2008) as well as http://www.quantise.maths-fi.com/

Theorem Dereich (2008)

 $\exists c > 0 \ \forall \mu, \sigma, x_0$

$$e^{\det}(n,\operatorname{Lip}(1)) \approx c \cdot \operatorname{E}\left(\int_0^T \sigma^2(X_t) \, dt\right)^{1/2} \cdot (\ln n)^{-1/2}.$$

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Question: construction of good quadrature formulas at reasonable cost?

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Question: construction of good quadrature formulas at reasonable cost? Formally,

- input: (μ, σ, x_0) as well as $n \in \mathbb{N}$,
- real number model with oracle for function/derivative values of μ and σ ,
- output: coefficients $a_i \in \mathbb{R}$ and nodes $x_i \in \mathfrak{X}$ (suitably coded).

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Theorem *Müller-Gronbach, R* (2011) $\forall \mu, \sigma, x_0$ construction of Q_n at cost O(n) such that

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Remark

- Analogous results for $\mathfrak{X} = L_p([0,T])$ instead of C([0,T]).
- Systems of SDEs require quantization of Lévy areas.
- Distance on the space of probability measures: Wasserstein metric.

$$dX_t = \alpha(\kappa - X_t) \, dt + \beta \sqrt{X_t} \, dW_t.$$

Implementation for $L_2([0,T])$ due to Toussaint (2008).

$$dX_t = \alpha(\kappa - X_t) \, dt + \beta \sqrt{X_t} \, dW_t.$$

1. For an equidistant time discretization, quantization of the marginal distribution

$$P_{(X_0, X_{T/m}, \dots, X_T)}$$

via quantized version of the Milstein scheme.



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 For an equidistant time discretization, quantization of the marginal distribution

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2. Local refinement via quantization of Brownian bridges, taking into account the local smoothness of X.

Basis functions

$$e_k(t) = \sqrt{2T} \cdot \sin(k\pi \cdot t/T)$$

for $k \in \mathbb{N}$.

Cf. adaptive step-size control for strong approximation, see Hofmann, Müller-Gronbach, R (2002), Müller-Gronbach (2002).

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Remark Alternative approaches to constructive quantization of SDEs

- ODE-based, using rough paths theory, see Luschgy, Pagès (2006), Pagès, Sellami (2010),
- using series expansions for X, see Luschgy, Pagès (2008).

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Remark Different techniques and results in the marginal case P_{X_T} , where

$$Q_n(f) = \sum_{i=1}^n a_i \cdot f(x_i)$$

with $a_i \in \mathbb{R}$ and $x_i \in \mathbb{R}^d$.

See Kusuoka (2001, 2004), Lyons, Victoir (2004), Crisan, Ghazali (2007), Litterer, Lyons (2010), Müller-Gronbach, R, Yaroslavtseva (2011), ... **Remark** Alternative approaches to constructive quantization of SDEs

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In 'Stochastic Computation' (B7)

S. Dereich: Constructive Quantization: approximation by empirical measures

L. Yaroslavtseva: A derandomization of the Euler scheme

Question: How to overcome the slow convergence of $e^{\det}(n, \operatorname{Lip}(1))$?

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 - better deterministic algorithms than just quadrature formulas?
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Question: Evaluation of integrands $f \in F$ anywhere in \mathfrak{X} at cost one?

III. Randomized Quadrature on the Lipschitz Class

Randomized quadrature formula $Q_n(f) = \sum_{i=1}^n a_i \cdot f(X_i)$.

Variable subspace sampling: for any scale of finite-dim. subspaces $\mathfrak{X}_0 \subset \mathfrak{X}_1 \subset \ldots \subset \mathfrak{X}$

$$X_1(\omega), \dots, X_n(\omega) \in \bigcup_{m=0}^{\infty} \mathfrak{X}_m,$$
$$\operatorname{cost}(Q_n) = \operatorname{E}\left(\sum_{i=1}^n \inf\{\dim \mathfrak{X}_m : X^{(i)} \in \mathfrak{X}_m\}\right).$$

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Example

 $\mathfrak{X}_m = \{x \in \mathfrak{X} \mid x \text{ piecewise linear with breakpoints } \ell/2^m \cdot T\}.$

Classical Euler-MC algorithm vs. multi-level Euler-MC algorithm.

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N-th minimal error, redefined,

$$e^{\operatorname{ran}}(N,F) = \inf_{\operatorname{cost}(Q) \le N} e(Q,F).$$

Theorem Creutzig, Dereich, Müller-Gronbach, R (2009)

$$N^{-1/2} \preceq e^{\operatorname{ran}}(N, \operatorname{Lip}(1)) \preceq N^{-1/2} \cdot \ln N.$$

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Remark

- Deterministic algorithms only yield $(\ln N)^{-1/2}$.
- Upper bound via multi-level Euler-MC algorithm.
- Fixed subspace sampling only yields $N^{-1/4}$, up to \ln 's.
- Lower bound valid for the class of randomized algorithms.

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Remark

- General result for
 - every probability measure P on any separable Banach space $\mathfrak X$ and

- F = Lip(1).

Upper and lower bounds for $e^{\operatorname{ran}}(N,F)$ in terms of

- quantization numbers of P,
- Kolmogorov widths of P, see *Mathé* (1990),

Motivation: As previously, $\mathfrak{X} = C([0,T])$,

$$dX_t = \mu(X_t) dt + \sigma(X_t) dW_t, \qquad t \in [0, T],$$
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and $f: \mathfrak{X} \to \mathbb{R}$. Approximate $I(f) = \int_{\mathfrak{X}} f \, dP_X$.

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and $f: \mathfrak{X} \to \mathbb{R}$. Approximate $I(f) = \int_{\mathfrak{X}} f \, dP_X$. Note that

 $X = \Gamma(\xi_1, \xi_2, \dots) =$ Euler expansion of X with step-sizes $1/2^k$,

based on Lévy-Ciesielski decomposition of W,

with ξ_1, ξ_2, \ldots iid and $P_{\xi_1} = N(0, 1) =: \rho$.

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Truncation

$$\int_{\mathbb{R}^d} f \circ \Gamma(z_1, \ldots, z_d, 0, \ldots) \, d\rho^{\otimes d}(z_1, \ldots, z_d).$$

The general formulation: Given

- a probability measure ρ on $D\subseteq \mathbb{R}$ and
- a class G of functions $g : \mathfrak{Z} \to \mathbb{R}$ on $\mathfrak{Z} = D^{\mathbb{N}}$.

Compute

$$I(g) = \int_{\mathfrak{Z}} g \, d\rho^{\otimes \mathbb{N}}.$$

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We study

• quadrature formulas

$$Q_n(g) = \sum_{i=1}^n a_i \cdot g(\mathbf{z}_i), \qquad a_i \in \mathbb{R}, \mathbf{z}_i \in \mathfrak{Z},$$

and variable subspace sampling,

• unit balls G in Hilbert spaces with a reproducing kernel.

Variable subspace sampling, based on

$$\mathfrak{Z}_{d,y} = \{ \mathbf{z} \in \mathfrak{Z} \mid z_{d+1} = z_{d+2} = \cdots = y \}$$

for any $y \in D$. Thus

$$\mathbf{z}_1,\ldots,\mathbf{z}_n\inigcup_{d=1}^\infty\mathfrak{Z}_{d,y}$$

and

$$\operatorname{cost}(Q_n) = \sum_{i=1}^n \inf\{d \mid \mathbf{z}_i \in \mathfrak{Z}_{d,y}\}.$$

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 $N\mbox{-th}\xspace$ minimal error

$$e^{\det}(N,G) = \inf_{\operatorname{cost}(Q) \le N} e(Q,G).$$

$$K_{\gamma}(\mathbf{y}, \mathbf{z}) = \prod_{j=1}^{\infty} (1 + \gamma_j \cdot \min(y_j, z_j)), \quad \mathbf{y}, \mathbf{z} \in \mathfrak{Z},$$

for weights $\gamma_1 \geq \gamma_2 \geq \cdots > 0$ such that

$$\sum_{j=1}^{\infty} \gamma_j < \infty.$$

See Hickernell, Wang (2002).

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Counterpart for functions $g: D^d \to \mathbb{R}$ of finitely many variables

- if $\gamma_1 = \cdots = \gamma_d = 1$ then $H(K_{\gamma}) = W_2^{(1,\dots,1)}([0,1]^d)$,
- for the weighted case see Sloan, Woźniakowski (1998), ..., Novak, Woźniakowski (2008, 2010, ...).

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$$\sum_{j=1}^{\infty} \gamma_j < \infty.$$

Example For

$$g(\mathbf{z}) = \sum_{j=1}^{\infty} \eta_j \cdot z_j^2$$

we have

$$g \in H(K_{\gamma}) \quad \Leftrightarrow \quad \sum_{j=1}^{\infty} \frac{\eta_j^2}{\gamma_j} < \infty.$$

$$K_{\gamma}(\mathbf{y}, \mathbf{z}) = \prod_{j=1}^{\infty} (1 + \gamma_j \cdot \min(y_j, z_j)), \qquad \mathbf{y}, \mathbf{z} \in \mathfrak{Z},$$

for weights $\gamma_1 \geq \gamma_2 \geq \cdots > 0$ such that

$$\sum_{j=1}^{\infty} \gamma_j < \infty.$$

Class of integrands

$$G_{\gamma} = \{ g \in H(K_{\gamma}) \mid ||g||_{\gamma} \le 1 \}.$$

Theorem Kuo, Sloan, Wasilkowski, Woźniakowski (2010)

Hickernell, Müller-Gronbach, Niu, R (2011), Gnewuch (2011),

Plaskota, Wasilkowski (2011)

Assume that ρ is the uniform distribution on D=[0,1] and

$$\gamma_j \asymp j^{-(1+2q)}$$

with q > 0.

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Assume that ρ is the uniform distribution on D = [0, 1] and

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$$N^{-\min(q,1)} \preceq e^{\det}(N, G_{\gamma}) \preceq N^{-\min(q,1)+\varepsilon}$$

 $\text{ if } |q-1| \geq 1/2 \text{ and }$

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Remark General results for product measures and weighted tensor product spaces.

Order of convergence of minimal errors in terms of decay of weights



Order of convergence of minimal errors in terms of decay of weights



Remark

- Variable subspace sampling superior to fixed subspace sampling.
- $W_2^1([0,1]) \hookrightarrow H(K_{\gamma})$, and $e^{\det}(N) \asymp N^{-1}$ on $W_2^1([0,1])$.
- Similar results for randomized (Monte Carlo) algorithms.

Order of convergence of minimal errors in terms of decay of weights



Remark

• Much stronger results concerning the computational cost in *Plaskota, Wasilkowski* (2011). A proof of the upper bound for $e^{\det}(N,G)$

• Construction of a multi-level algorithm

Put $g_d(\mathbf{z}) = g(z_1, ..., z_d, 0, ...)$. Clearly

$$g_{d_L}(\mathbf{z}) = g_{d_1}(\mathbf{z}) + \sum_{\ell=2}^{L} (g_{d_\ell}(\mathbf{z}) - g_{d_{\ell-1}}(\mathbf{z})).$$

Use rank-1 lattice rules for integration of g_{d_1} and $g_{d_\ell} - g_{d_{\ell-1}}$.

A proof of the upper bound for $e^{\det}(N,G)$

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Analysis

Choose suitable weights $\widetilde{\gamma}_j \gg \gamma_j$, derive estimates for

$$\sup_{\|g\|_{\boldsymbol{\gamma}}\leq 1}\|g_{d_{\ell}}-g_{d_{\ell-1}}\|_{\boldsymbol{\widetilde{\gamma}}}.$$

Employ tractability result for rank-1 lattice rules, see *Hickernell, Sloan, Wasilkowski* (2004).

A Numerical Example

Black-Scholes Model

$$dX_t = \alpha \cdot X_t \, dt + \beta \cdot X_t \, dW_t,$$
$$X_0 = x_0$$

with $t \in [0, 1]$, $\alpha = 0.05$, $\beta = 0.5$, and $x_0 = 2.0$. Asian option

$$f(X) = \left(\int_0^1 X(t) \, dt - K\right)_+$$

with K = 2.0. See Giles, Waterhouse (2009).

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Estimation of the root mean square error based on 1000 replications, see Mayer (2011). Log-linear regression, assuming that

rms error =
$$c \cdot \frac{(\ln(\cot))^{\beta_1}}{\cot^{\beta_2}}$$
.

Root mean square error vs. cost



Root mean square error vs. cost



Asymptotic results



Root mean square error vs. cost



Lévy-Ciesielski, shifted rank-one ?

Lévy-Ciesielski, shifted rank-one, multi-level ?